

# Broadcast-Optimal Two Round MPC with Asynchronous Peer-to-Peer Channels

*Ivan Damgård<sup>1</sup>, Divya Ravi<sup>1,2</sup>, Luisa Siniscalchi<sup>3</sup>, Sophia Yakoubov<sup>1</sup>*

[eprint.iacr.org/2023/1187](https://eprint.iacr.org/2023/1187)

<sup>1</sup>Aarhus University, <sup>2</sup>University of Amsterdam, <sup>3</sup>DTU

# Broadcast-Optimal Two Round MPC

# Broadcast-Optimal Two Round MPC



**Grace**



**Bruce**



**Alan**

# Broadcast-Optimal Two Round MPC



**Grace**

$X_G$



**Bruce**

$X_B$



**Alan**

$X_A$

# Broadcast-Optimal Two Round MPC

$$y = f(x_A, x_B, x_G)$$



**Grace**

$x_G$



**Bruce**

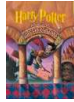
$x_B$

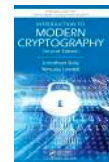
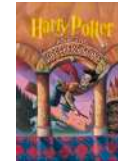


**Alan**

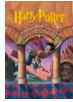
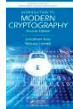
$x_A$

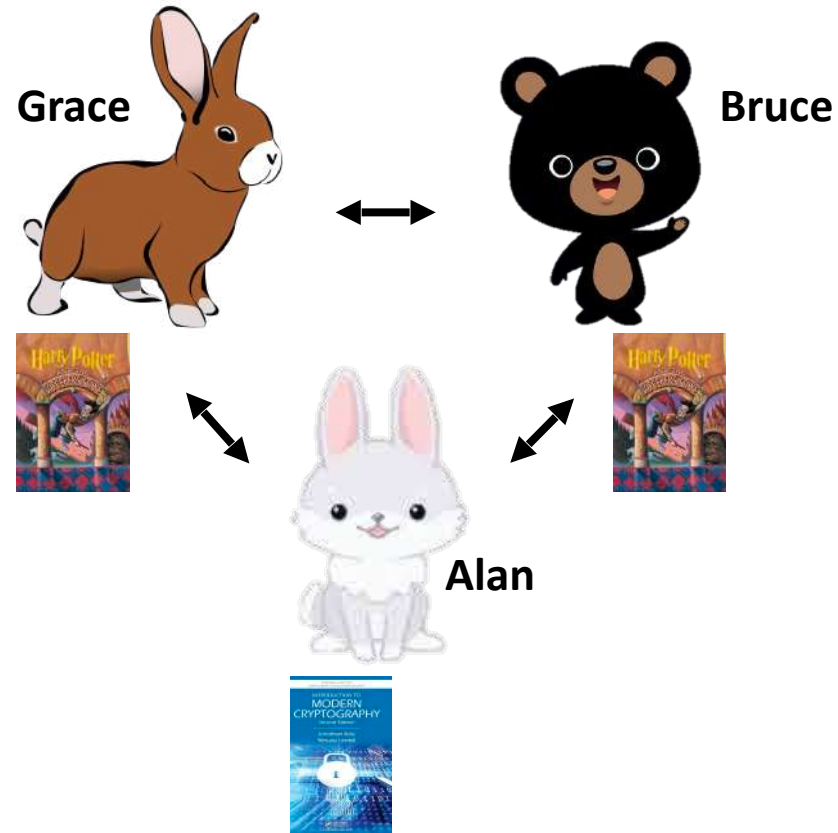
# Broadcast-Optimal Two Round MPC

y = should we read  or  for book club?

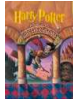
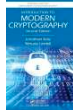


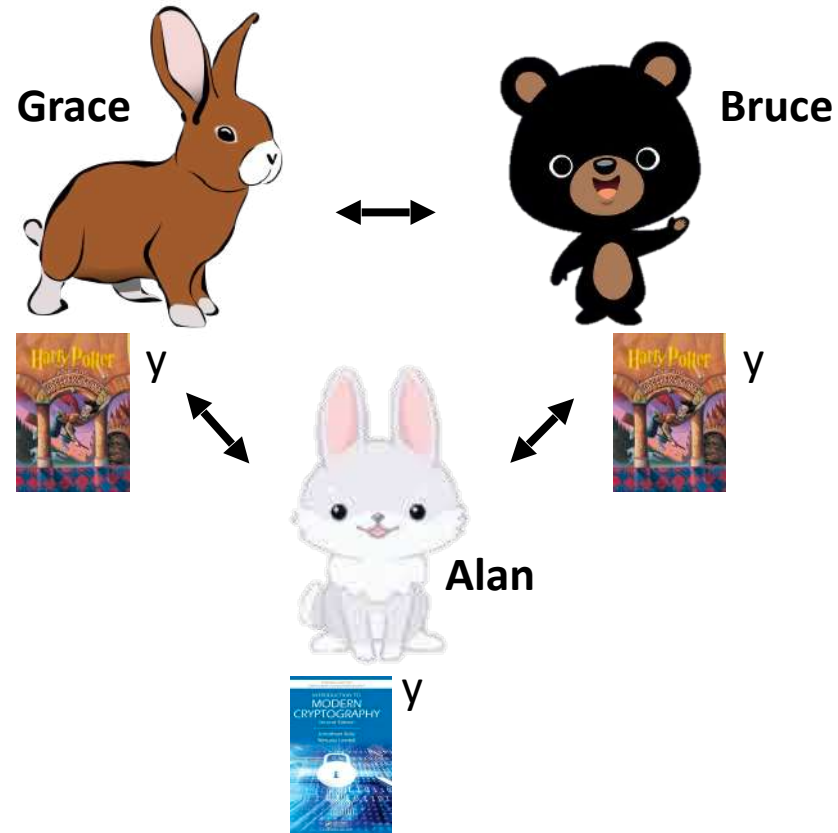
# Broadcast-Optimal Two Round MPC

y = should we read  or  for book club?



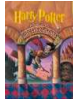
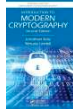
# Broadcast-Optimal Two Round MPC

y = should we read  or  for book club?

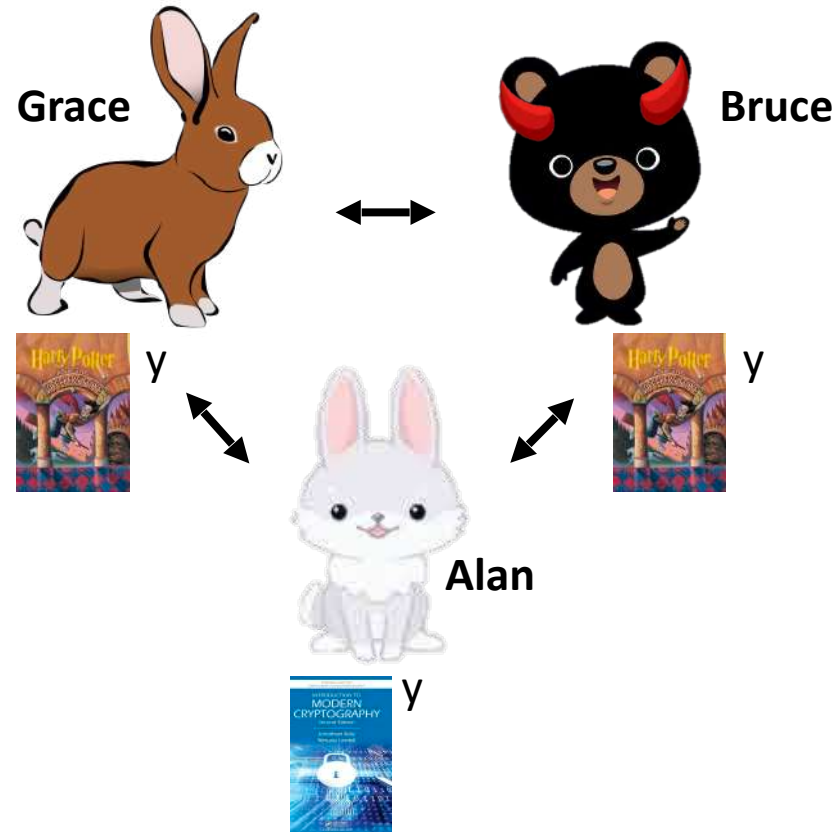




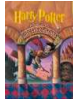
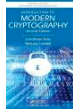
# Broadcast-Optimal Two Round MPC

$y =$  should we read  or  for book club?

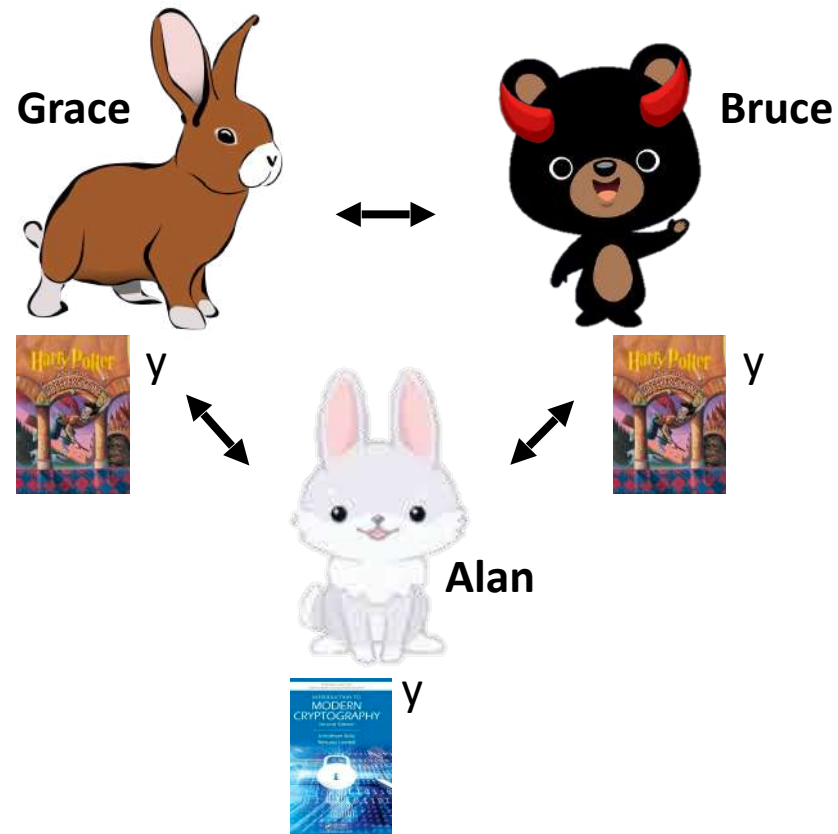
- Privacy:  $t$  corrupt parties learn no additional information about honest parties' inputs



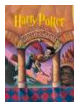
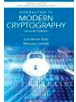
# Broadcast-Optimal Two Round MPC

$y =$  should we read  or  for book club?

- Privacy:  $t$  corrupt parties learn no additional information about honest parties' inputs
- Correctness

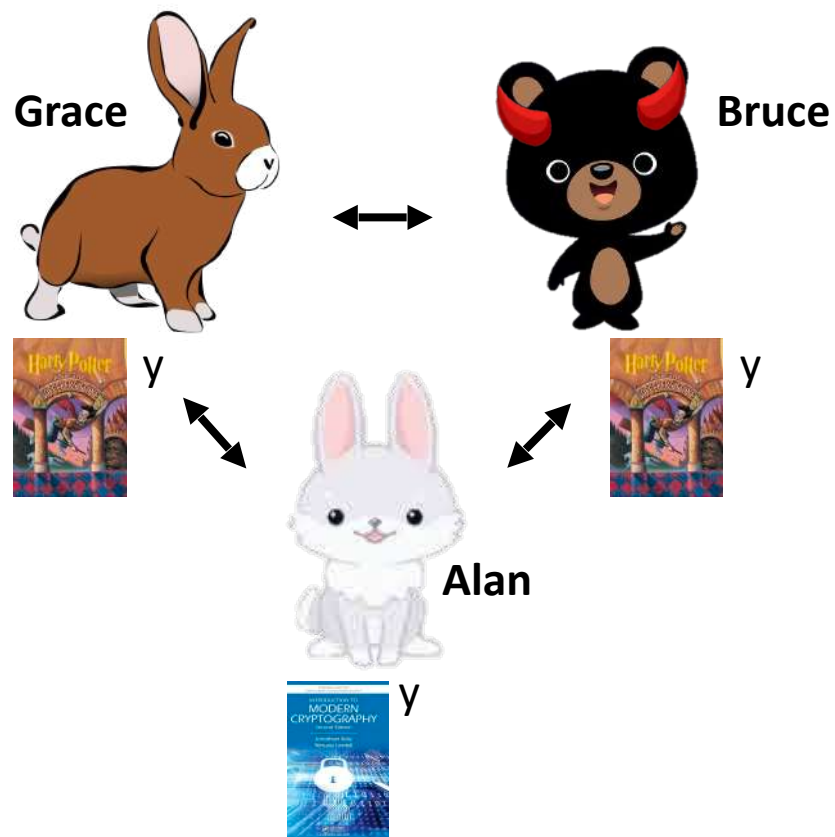


# Broadcast-Optimal Two Round MPC

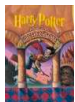
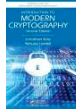
$y =$  should we read  or  for book club?

- Privacy:  $t$  corrupt parties learn no additional information about honest parties' inputs
- Correctness

Stronger  
↓

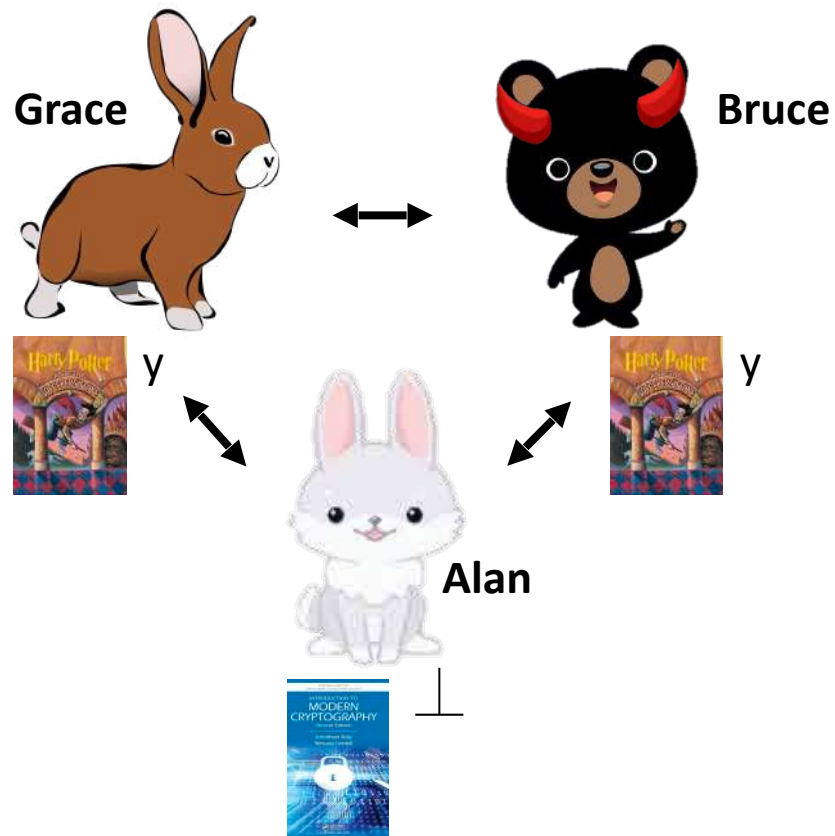


# Broadcast-Optimal Two Round MPC

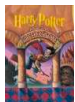
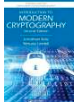
y = should we read  or  for book club?

- Privacy: t corrupt parties learn no additional information about honest parties' inputs
- Correctness
  - **Selective Abort**

Stronger ↓

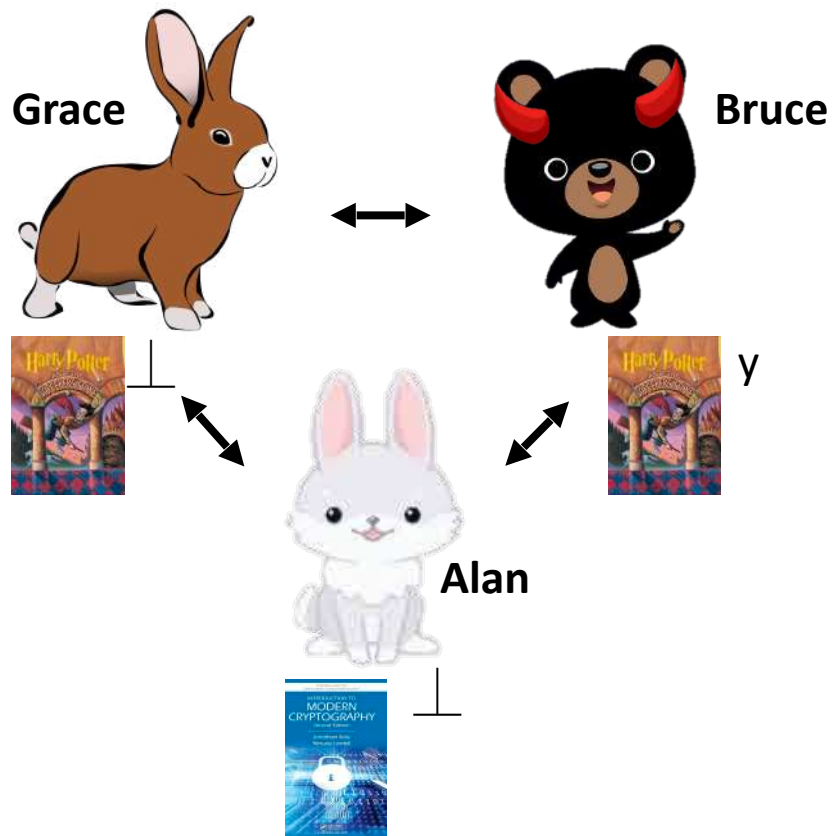


# Broadcast-Optimal Two Round MPC

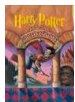
y = should we read  or  for book club?

- Privacy: t corrupt parties learn no additional information about honest parties' inputs
- Correctness
  - Selective Abort
  - **Unanimous Abort**

Stronger ↓

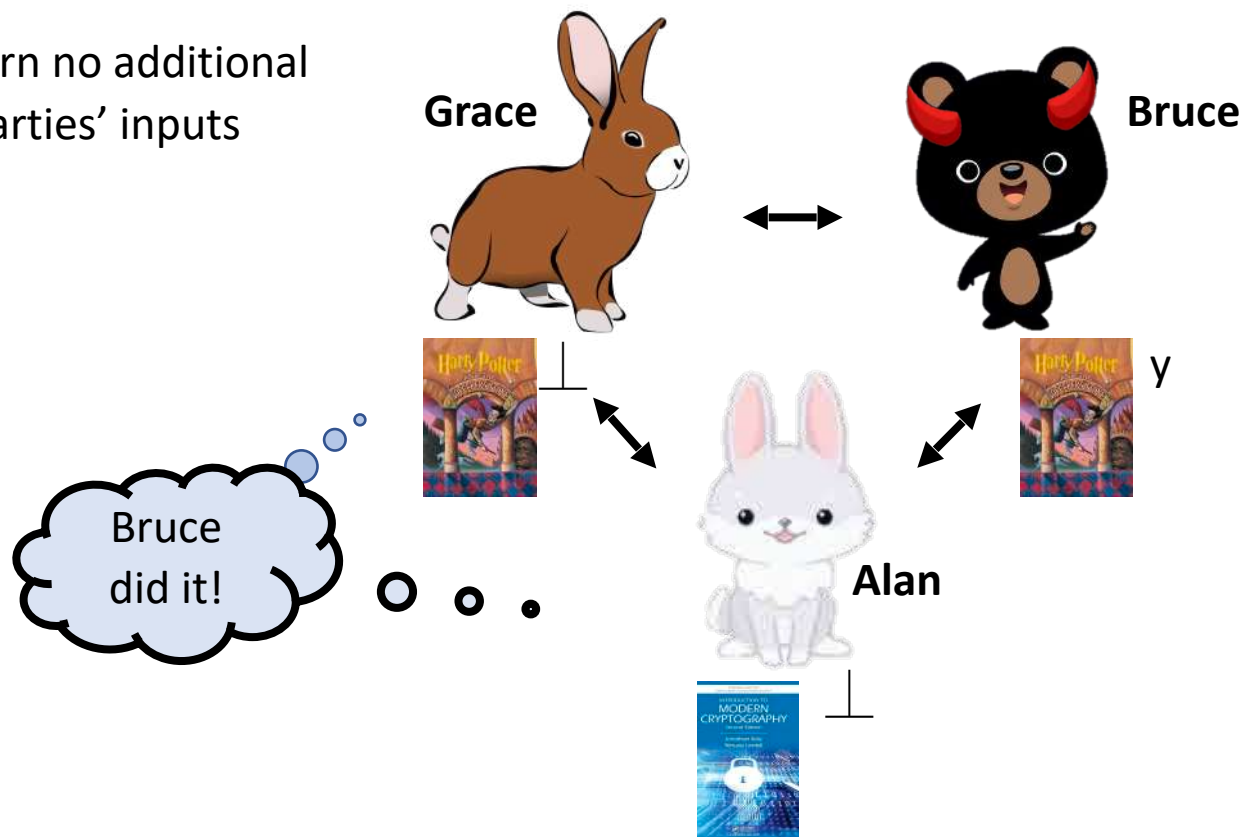


# Broadcast-Optimal Two Round MPC

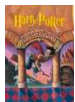
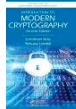
$y =$  should we read  or  for book club?

- Privacy:  $t$  corrupt parties learn no additional information about honest parties' inputs
- Correctness:
  - Selective Abort
  - Unanimous Abort
  - **Identifiable Abort**

Stronger  
↓

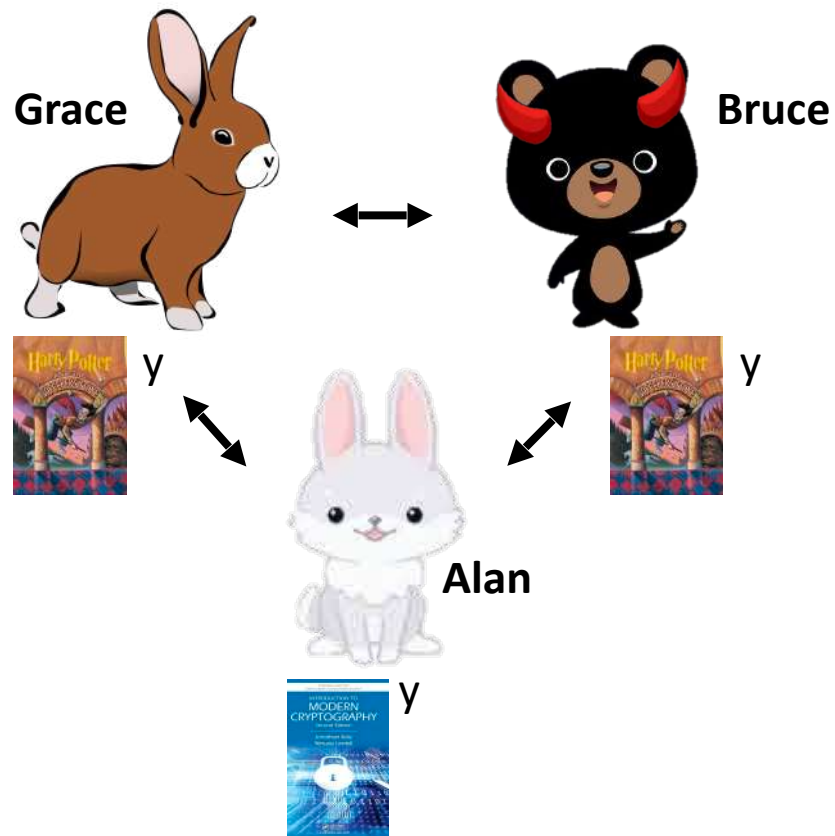


# Broadcast-Optimal Two Round MPC

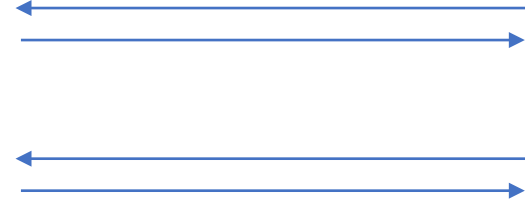
y = should we read  or  for book club?

- Privacy: t corrupt parties learn no additional information about honest parties' inputs
- Correctness:
  - Selective Abort
  - Unanimous Abort
  - Identifiable Abort
  - **Guaranteed Output Delivery**

Stronger ↓

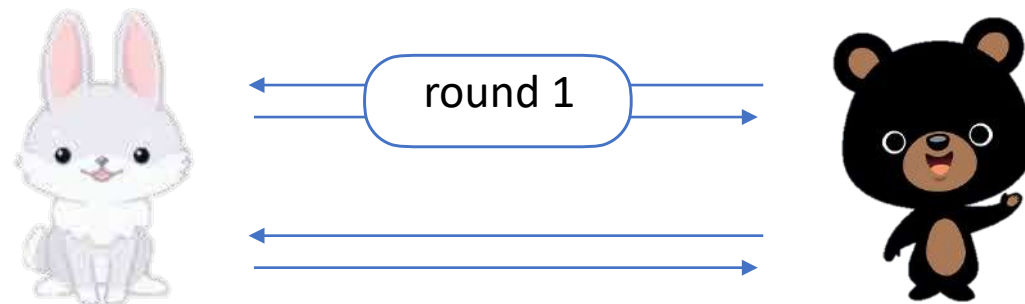


# Broadcast-Optimal Two Round MPC

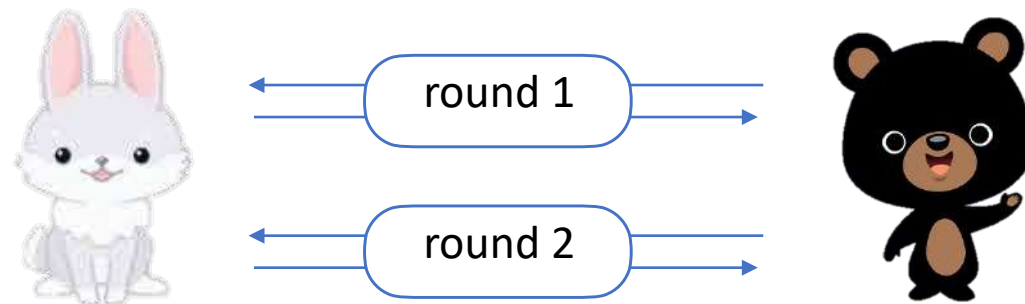




# Broadcast-Optimal Two Round MPC

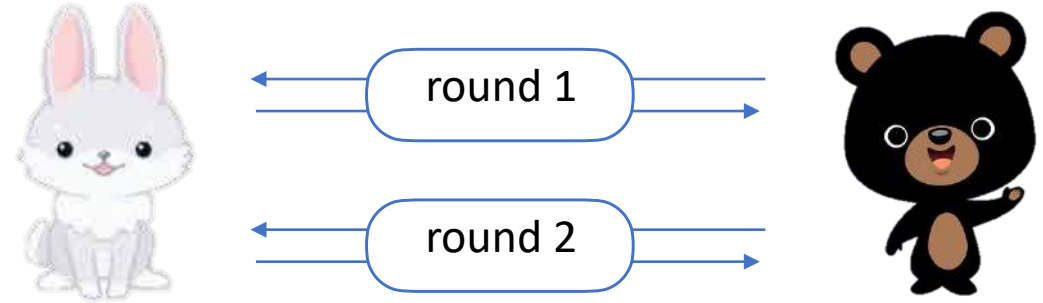


# Broadcast-Optimal Two Round MPC



# Broadcast-Optimal Two Round MPC

- Rounds are expensive!
  - At least two rounds needed for MPC



# Broadcast-Optimal Two Round MPC

- Rounds are expensive!
  - At least two rounds needed for MPC
- Broadcast is expensive!
  - Takes many rounds, or
  - Uses expensive resources



# Broadcast-Optimal Two Round MPC

- Rounds are expensive!
  - At least two rounds needed for MPC
- Broadcast is expensive!
  - Takes many rounds, or
  - Uses expensive resources
- Most two-round MPC:
  - Does not use broadcast (gets weaker guarantees - e.g. selective abort)

# Broadcast-Optimal Two Round MPC

- Rounds are expensive!
  - At least two rounds needed for MPC
- Broadcast is expensive!
  - Takes many rounds, or
  - Uses expensive resources
- Most two-round MPC:
  - Does not use broadcast (gets weaker guarantees - e.g. selective abort)
  - Uses broadcast in both rounds (expensive)

# Broadcast-Optimal Two Round MPC

which of the two rounds do we *really* need broadcast in?

# Broadcast-Optimal Two Round MPC

which of the two rounds do we *really* need broadcast in?

	Dishonest majority	Honest majority
PKI	[CGZ20]	[DMRSY21]
No PKI		[DRSY23]

PKI: public key infrastructure



# Broadcast-Optimal Two Round MPC

which of the two rounds do we *really* need broadcast in?

	Dishonest majority	Honest majority
PKI	[CG2201]	[DRSY23]
No PKI	<i>Assumes Synchronous Channels</i>	[DRSY23]

PKI: public key infrastructure

# Synchronous communication



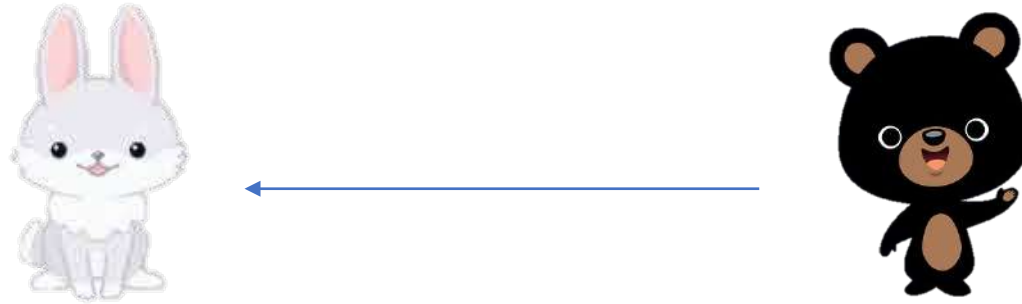
# Synchronous communication

Guaranteed to be delivered within one round

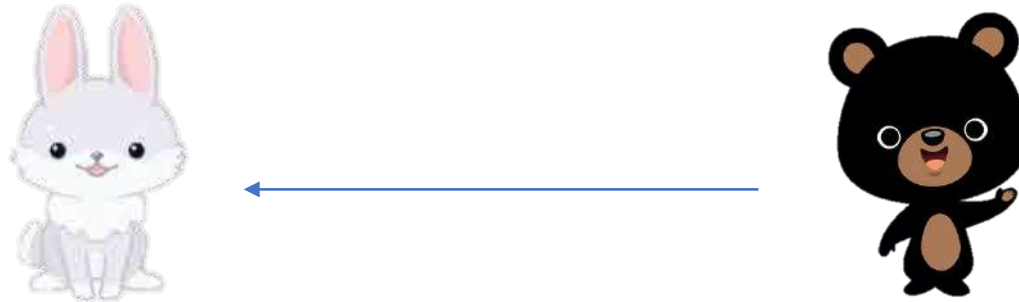


Guaranteed to be delivered within one round

Synchronous  
communication



Asynchronous  
communication



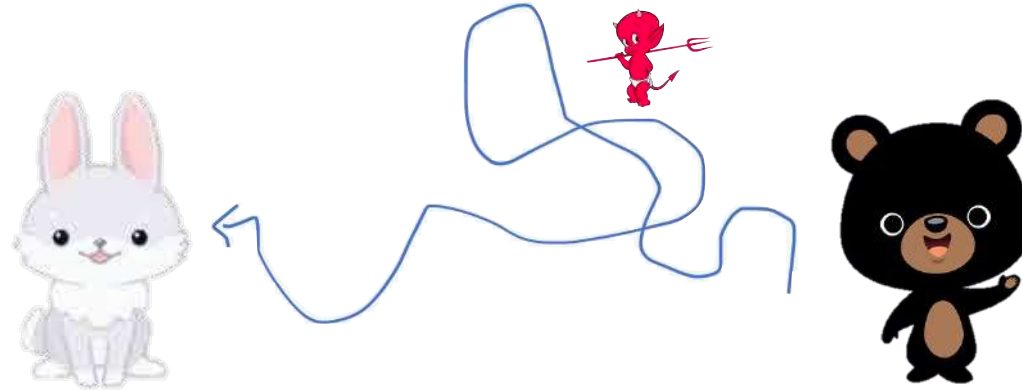
# Synchronous communication

Guaranteed to be delivered within one round



---

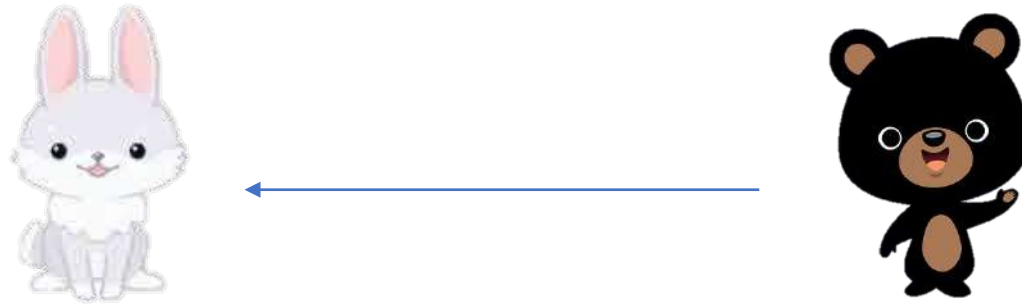
# Asynchronous communication



Arbitrarily delayed by adversary

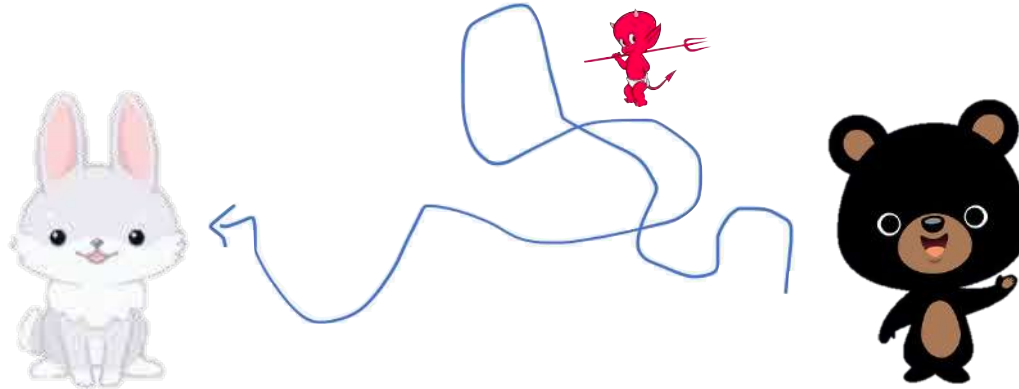
# Synchronous communication

Guaranteed to be delivered within one round



---

# Asynchronous communication



Arbitrarily delayed by adversary

Honest parties never know whether...

- message was never sent, or
- It was delayed

# Broadcast-Optimal Two Round MPC with Asynchronous Peer-to-Peer Rounds

Second round

		asynchronous P2P	BC
first round	asynchronous P2P	impossible with standard definitions of security	
	BC		well-studied

# Broadcast-Optimal Two Round MPC with Asynchronous Peer-to-Peer Rounds

Second round

		asynchronous P2P	BC
first round	asynchronous P2P	impossible with standard definitions of security	
	BC	Impossible for $n \leq 2t$ Possible otherwise (under some conditions)	well-studied



# Broadcast-Optimal Two Round MPC with Asynchronous Peer-to-Peer Rounds

Second round

		Second round	
		asynchronous P2P	BC
first round	asynchronous P2P	impossible with standard definitions of security	Impossible with classical notion of asynchrony We introduce a new variant!
	BC	Impossible for $n \leq 2t$ Possible otherwise (under some conditions)	well-studied

# Broadcast-Optimal Two Round MPC with Asynchronous Peer-to-Peer Rounds

Second round

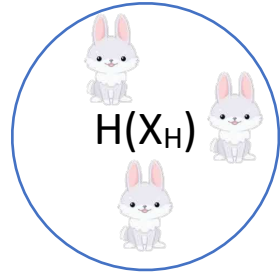
		asynchronous P2P	BC
first round	asynchronous P2P	impossible with standard definitions of security	Impossible with classical notion of asynchrony We introduce a new variant!
	BC	Impossible for $n \leq 2t$ Possible otherwise (under some conditions)	well-studied

# Impossibility of asyncP2P, BC

# Impossibility of asyncP2P, BC



# Impossibility of asyncP2P, BC



# Impossibility of asyncP2P, BC

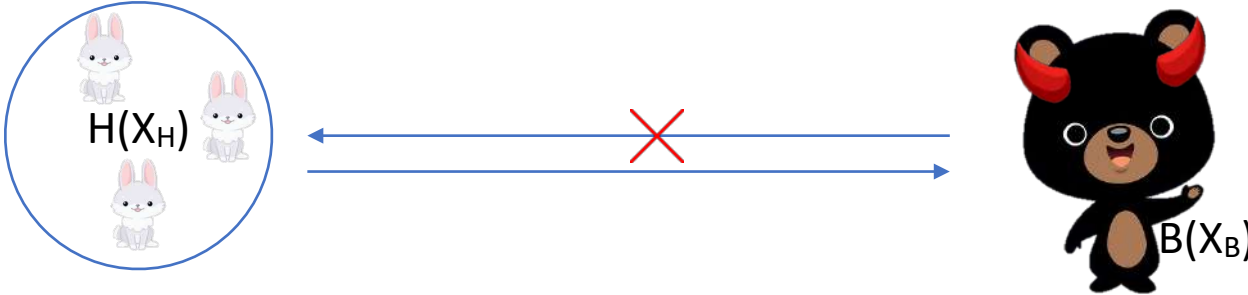


round 1: asyncP2P

round 2: BC



# Impossibility of asyncP2P, BC



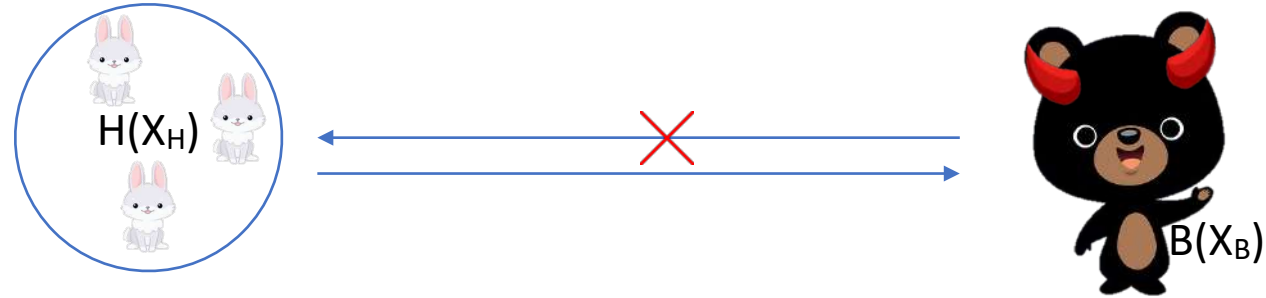
round 1: asyncP2P

round 2: BC



- Nothing H says depends on  $X_B$

# Impossibility of asyncP2P, BC



round 1: asyncP2P

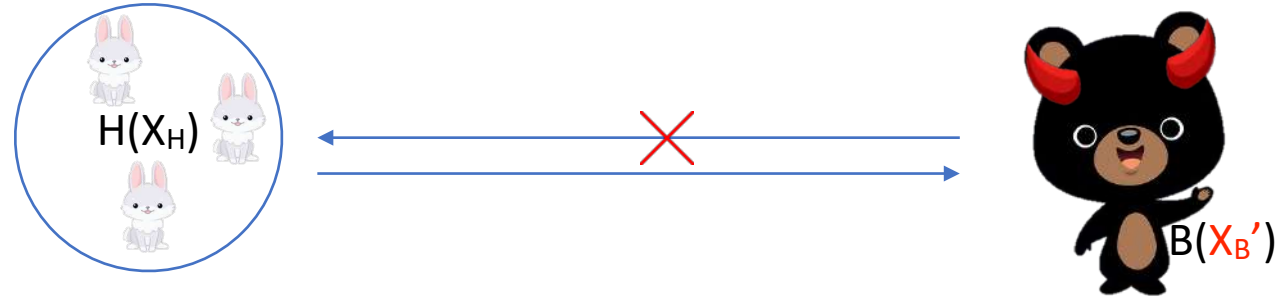
round 2: BC





# Impossibility of asyncP2P, BC

- Nothing H says depends on  $X_B$
- ... so B can swap out  $X_B$  *after the computation!*



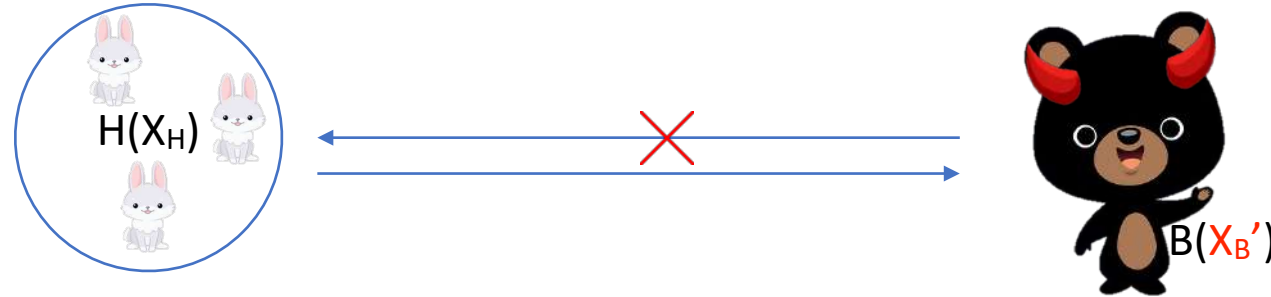
round 1: asyncP2P

round 2: BC



# Impossibility of asyncP2P, BC

- Nothing H says depends on  $X_B$
- ... so B can swap out  $X_B$  *after the computation!*



round 1: asyncP2P

round 2: BC



$(t_d, t_m)$ -asynchrony

# $(t_d, t_m)$ -asynchrony

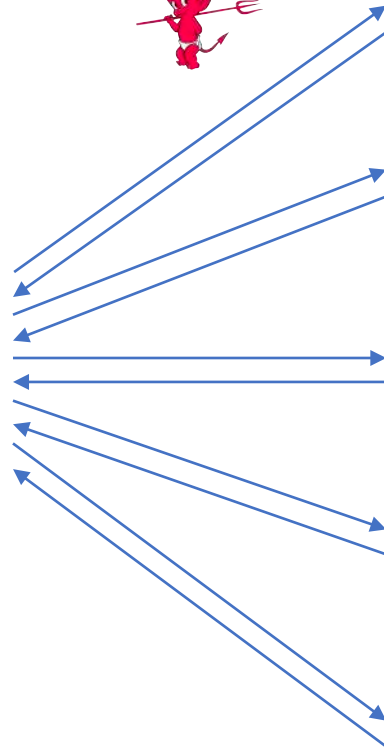
deafness  
threshold

muteness  
threshold

# $(t_d, t_m)$ -asynchrony

deafness  
threshold

muteness  
threshold

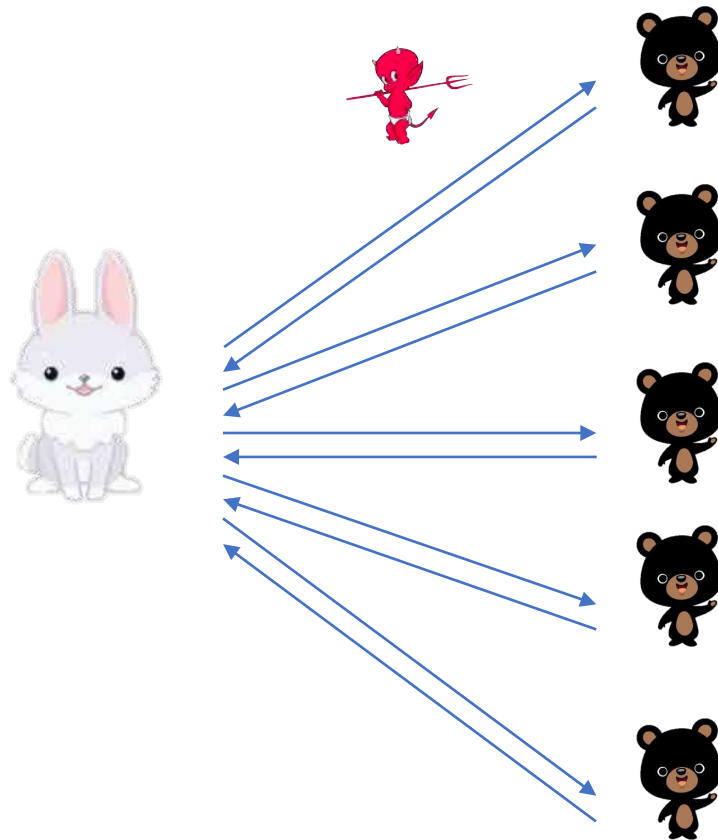


# $(t_d, t_m)$ -asynchrony

deafness  
threshold

muteness  
threshold

- $t_d = 2$

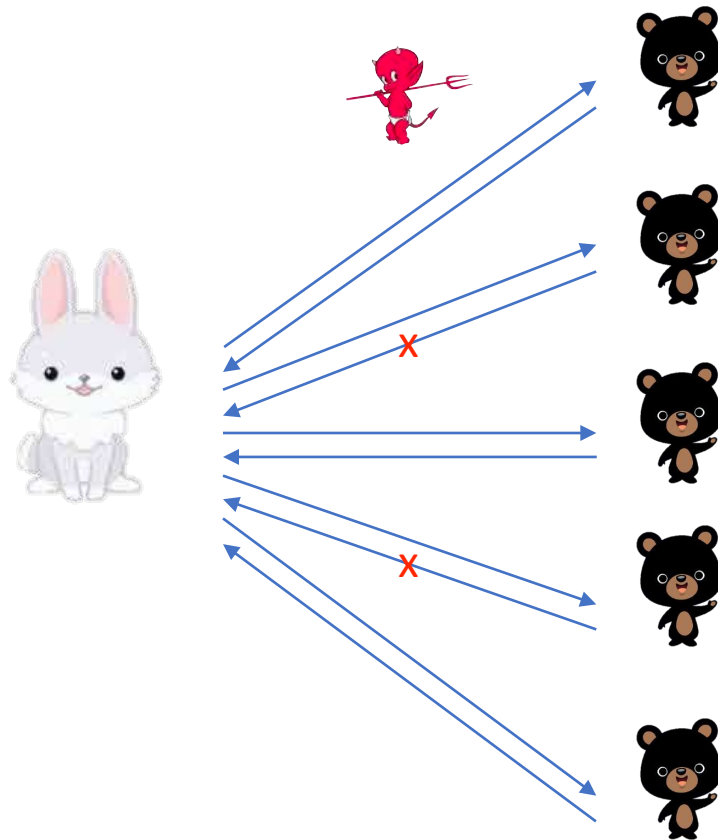


# $(t_d, t_m)$ -asynchrony

deafness  
threshold

muteness  
threshold

- $t_d = 2$

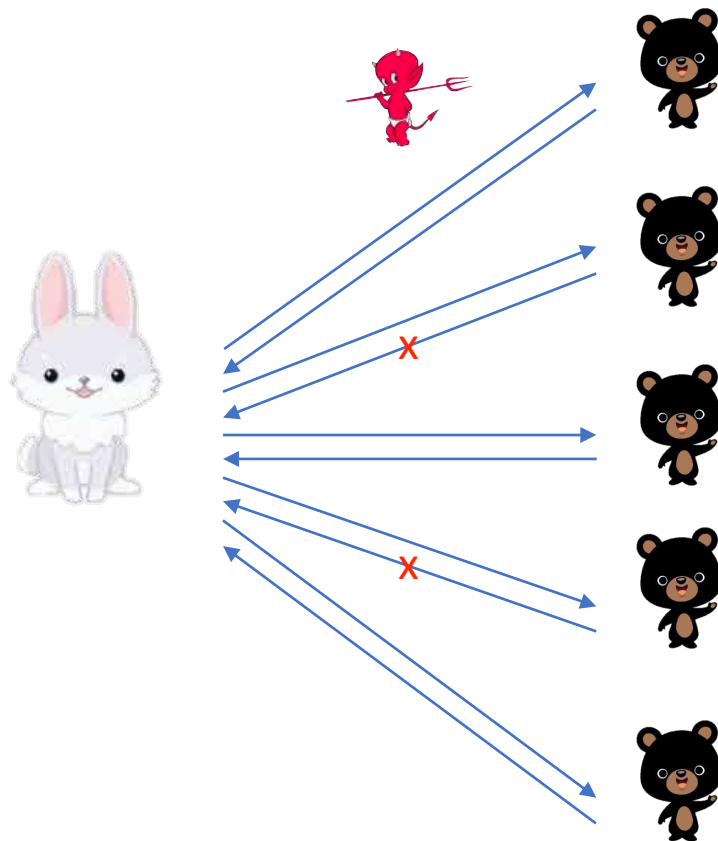


# $(t_d, t_m)$ -asynchrony

deafness  
threshold

muteness  
threshold

- $t_d = 2$
- $t_m = 1$



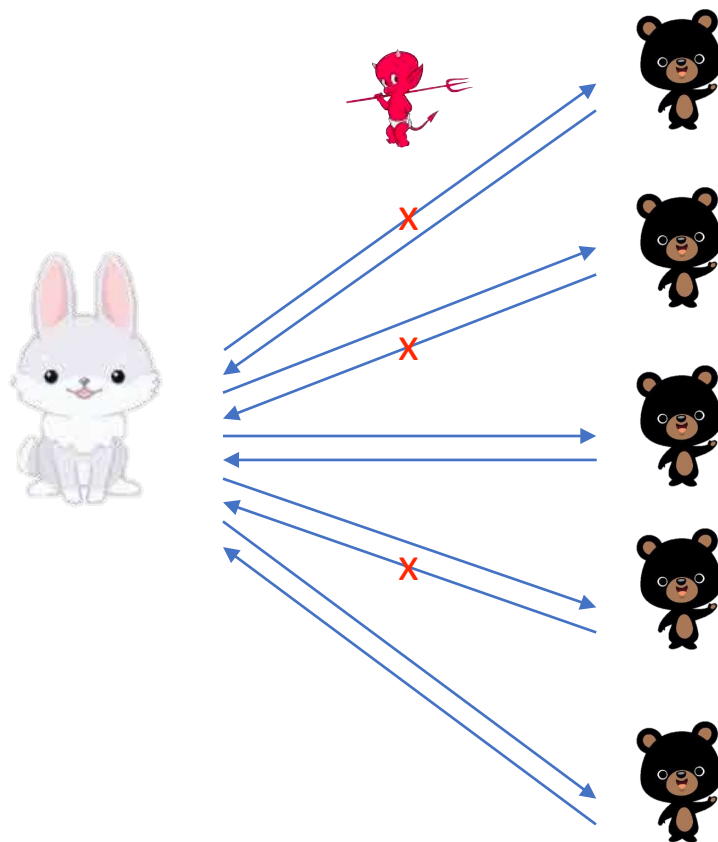


# $(t_d, t_m)$ -asynchrony

deafness  
threshold

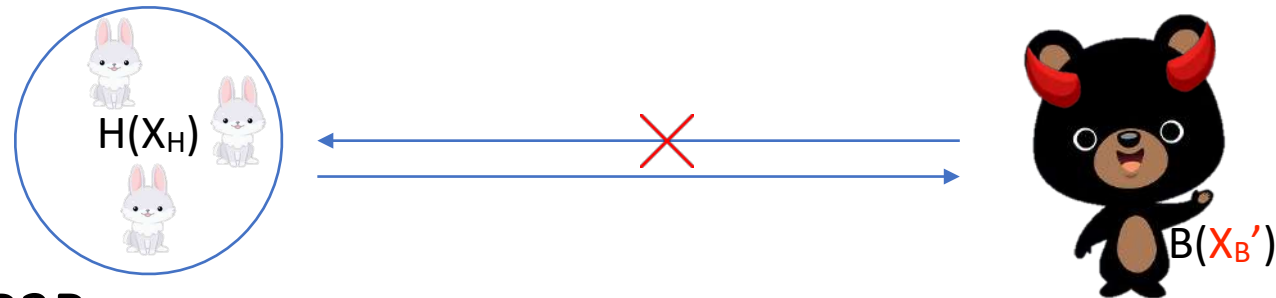
muteness  
threshold

- $t_d = 2$
- $t_m = 1$



# $(t_d, t_m)$ -asyncP2P, BC

with PKI



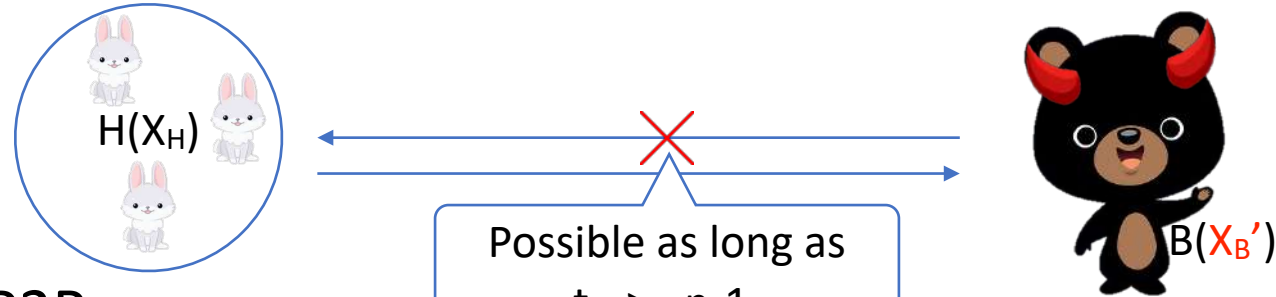
round 1:  $(t_d, t_m)$ -asyncP2P

round 2: BC



# $(t_d, t_m)$ -asyncP2P, BC

with PKI

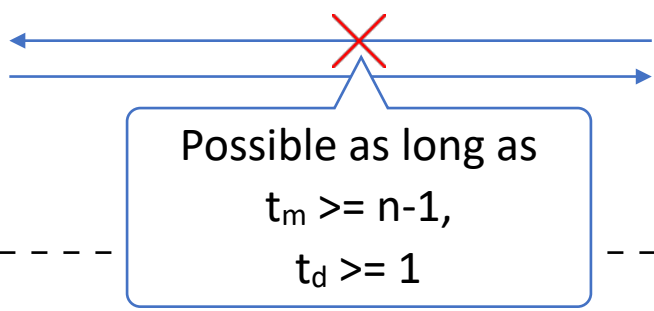
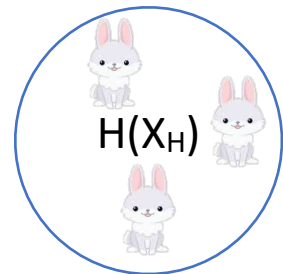
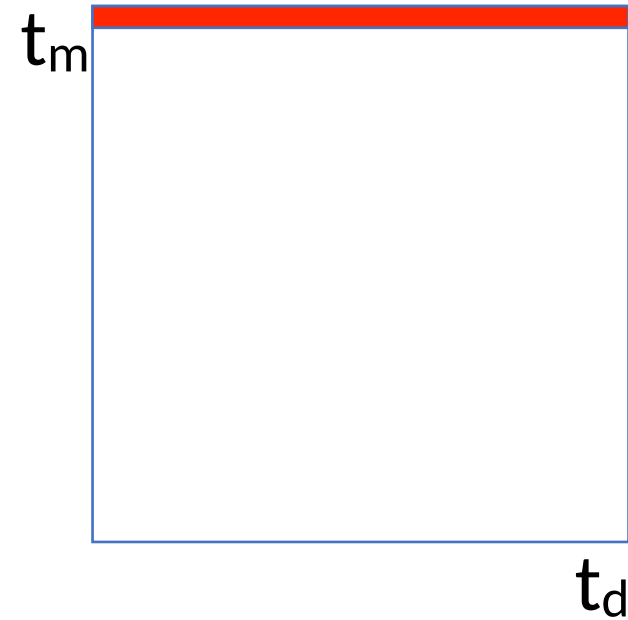


round 2: BC

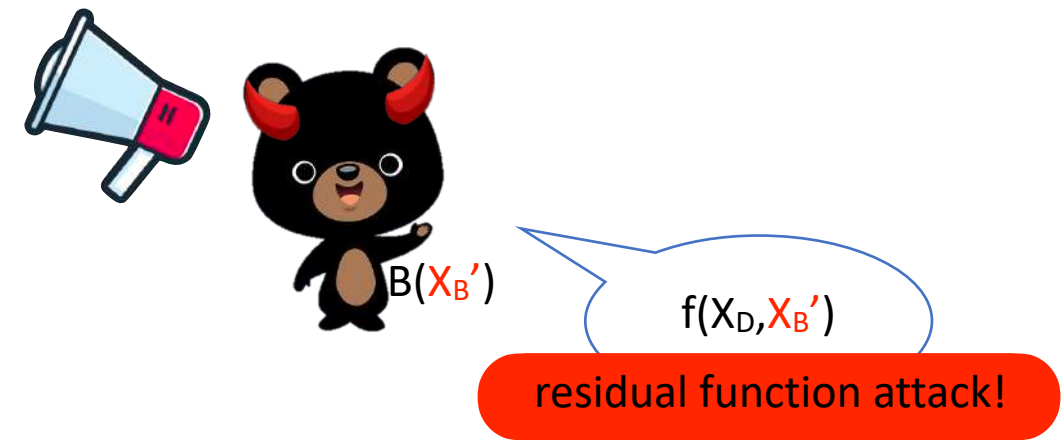
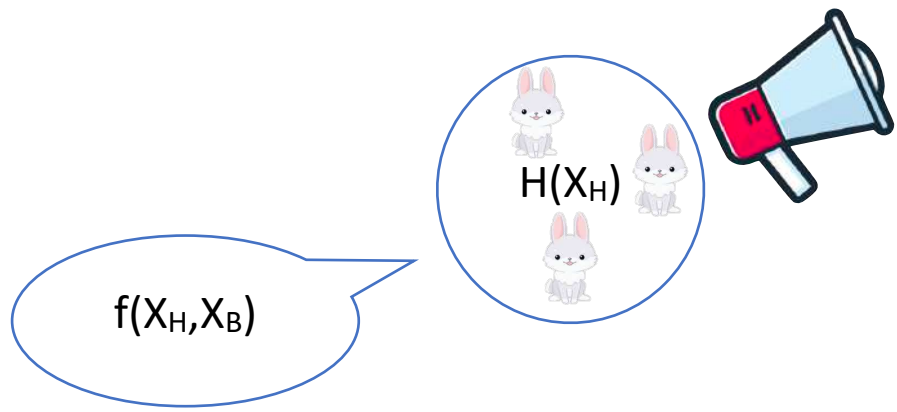


# $(t_d, t_m)$ -asyncP2P, BC

with PKI

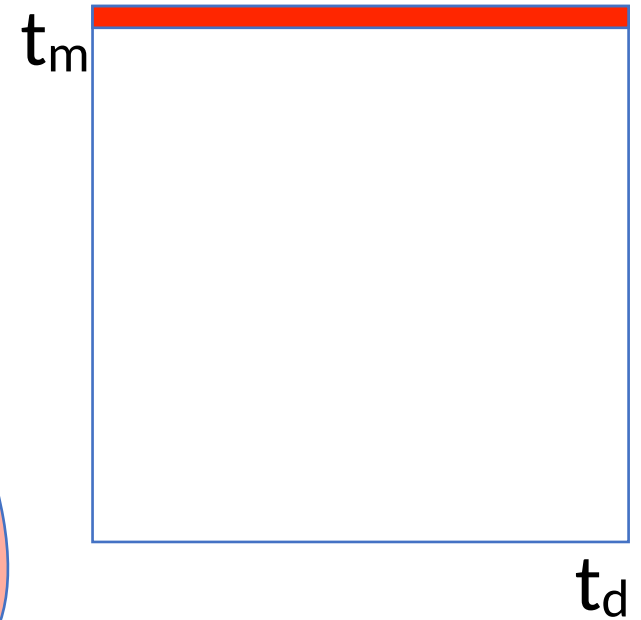
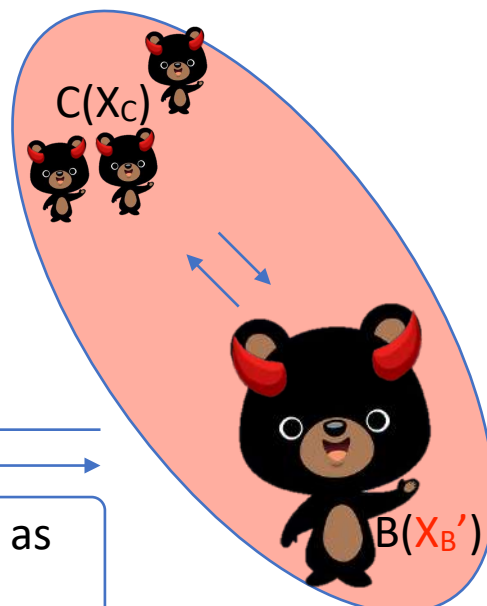
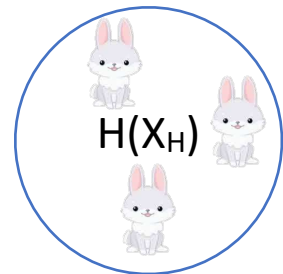


round 1:  $(t_d, t_m)$ -asyncP2P  
-----  
round 2: BC



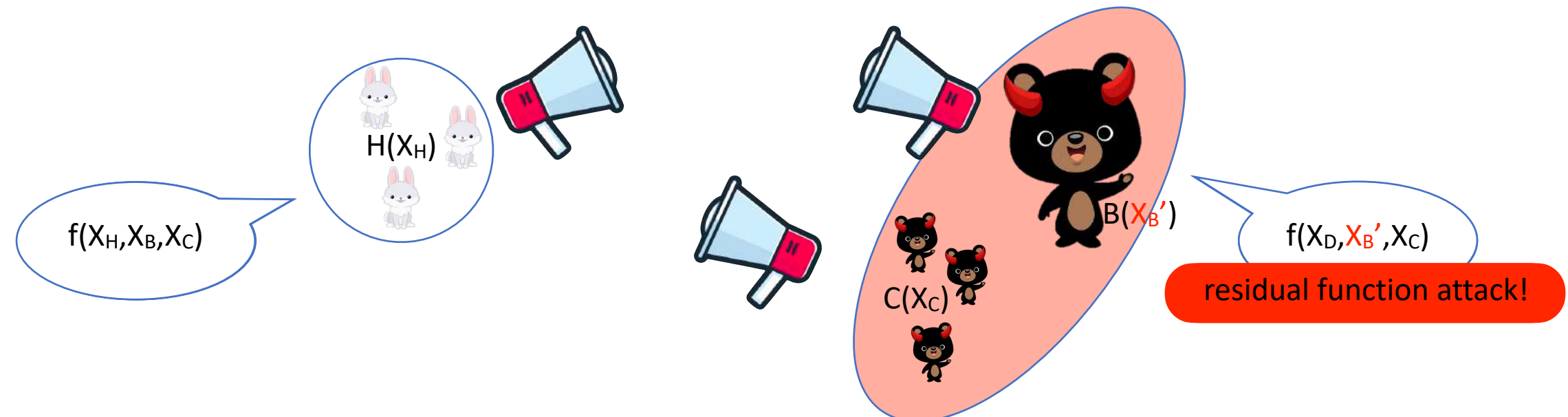
# $(t_d, t_m)$ -asyncP2P, BC

with PKI



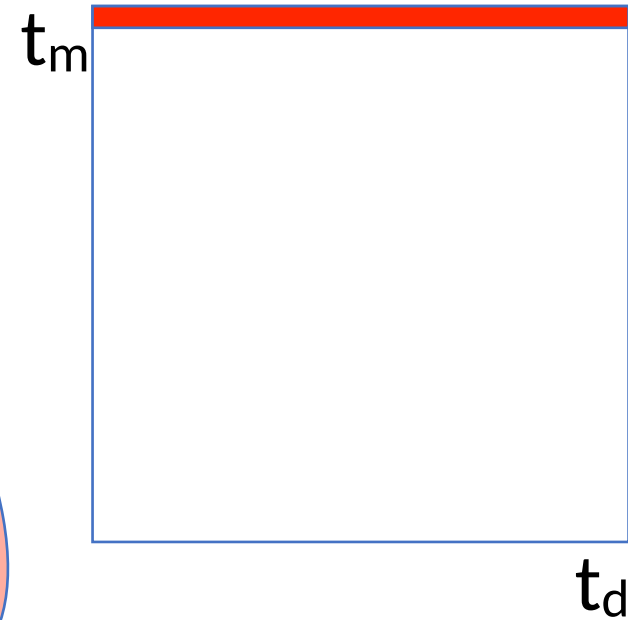
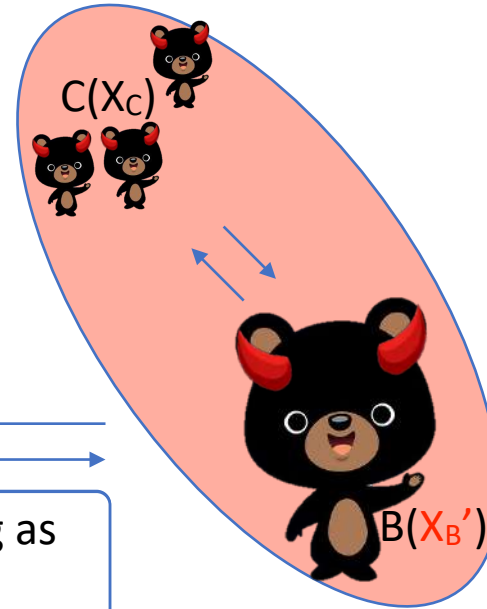
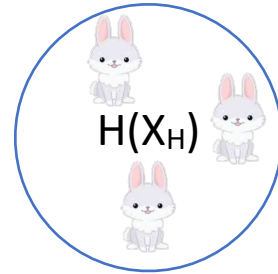
Possible as long as  
 $t_m \geq n-1,$   
 $t_d \geq 1$

round 1:  $(t_d, t_m)$ -asyncP2P  
-----  
round 2: BC



# $(t_d, t_m)$ -asyncP2P, BC

with PKI



Possible as long as  
 $t_m \geq n-t,$   
 $t_d \geq 1$

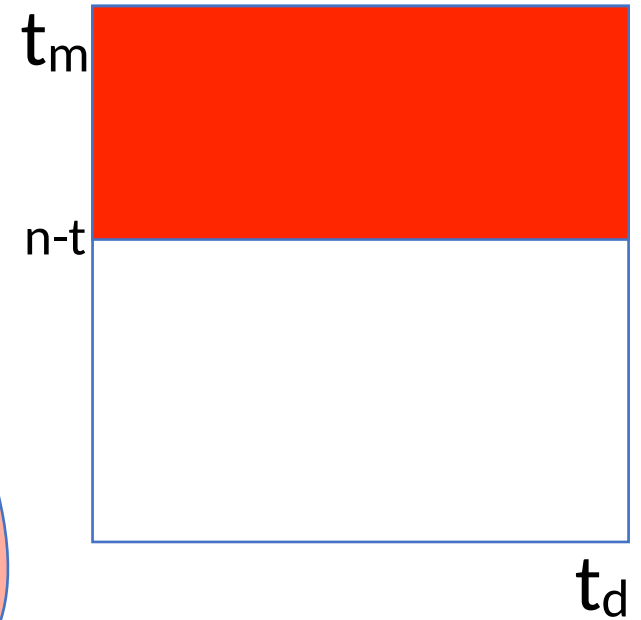
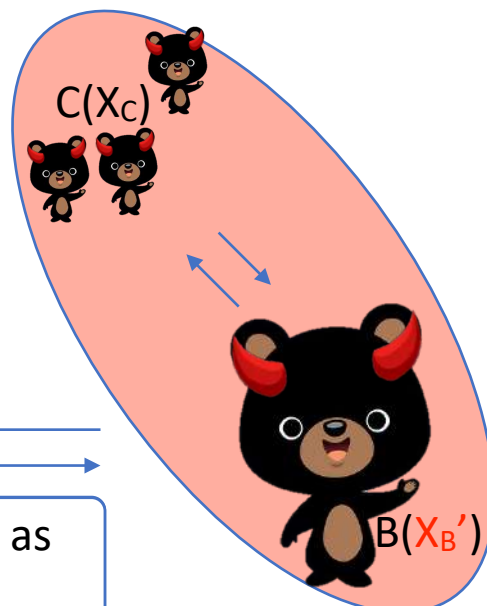
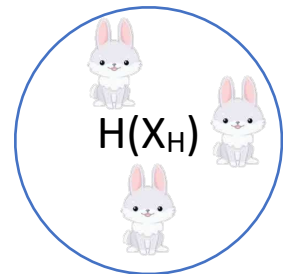
round 1:  $(t_d, t_m)$ -asyncP2P

round 2: BC



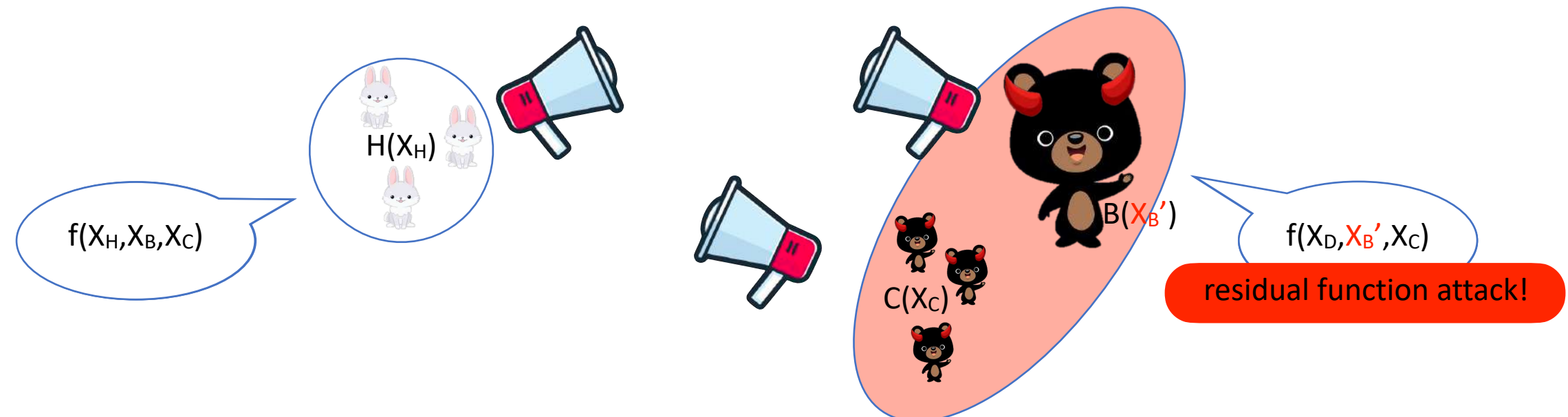
# $(t_d, t_m)$ -asyncP2P, BC

with PKI



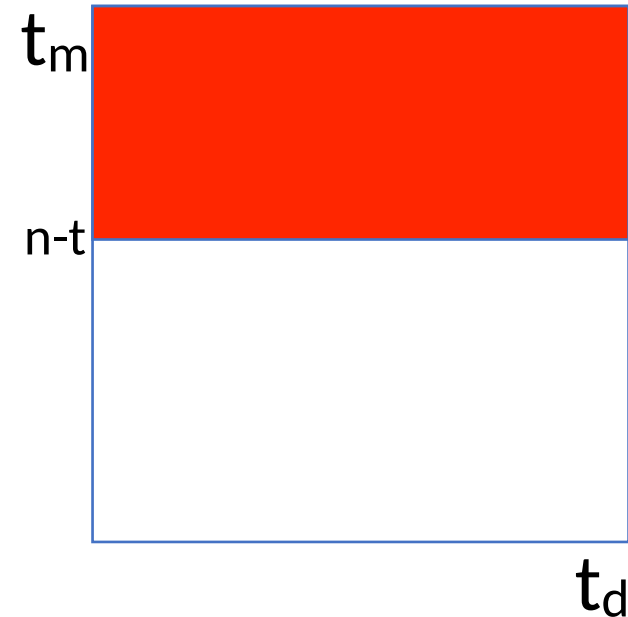
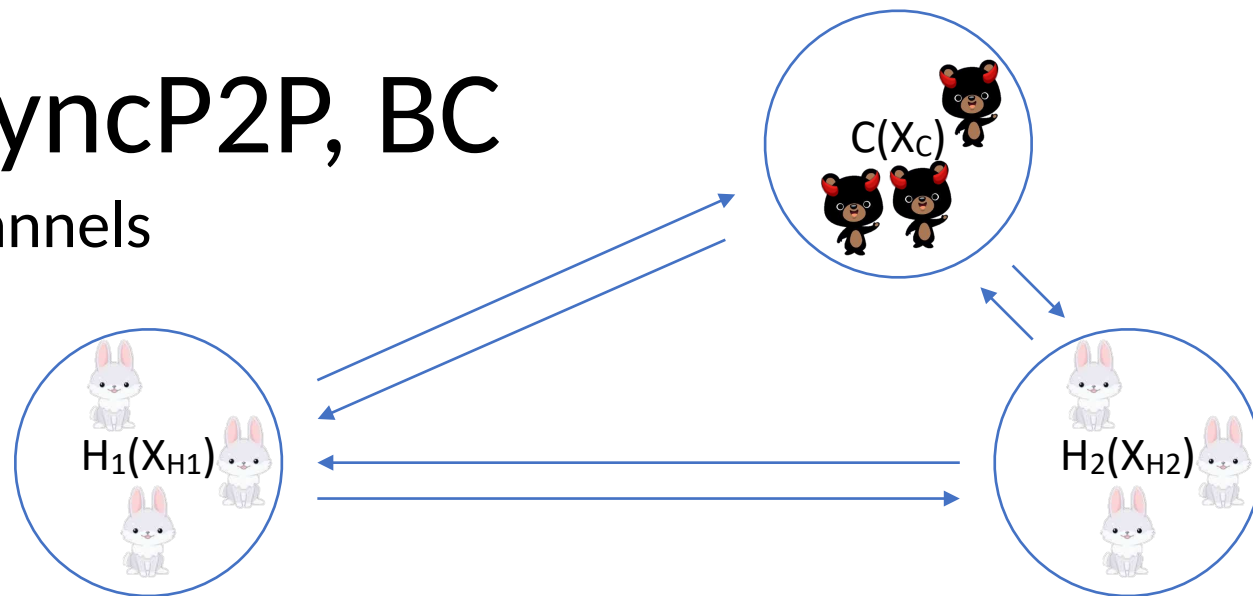
Possible as long as  
 $t_m \geq n-t,$   
 $t_d \geq 1$

round 1:  $(t_d, t_m)$ -asyncP2P  
-----  
round 2: BC



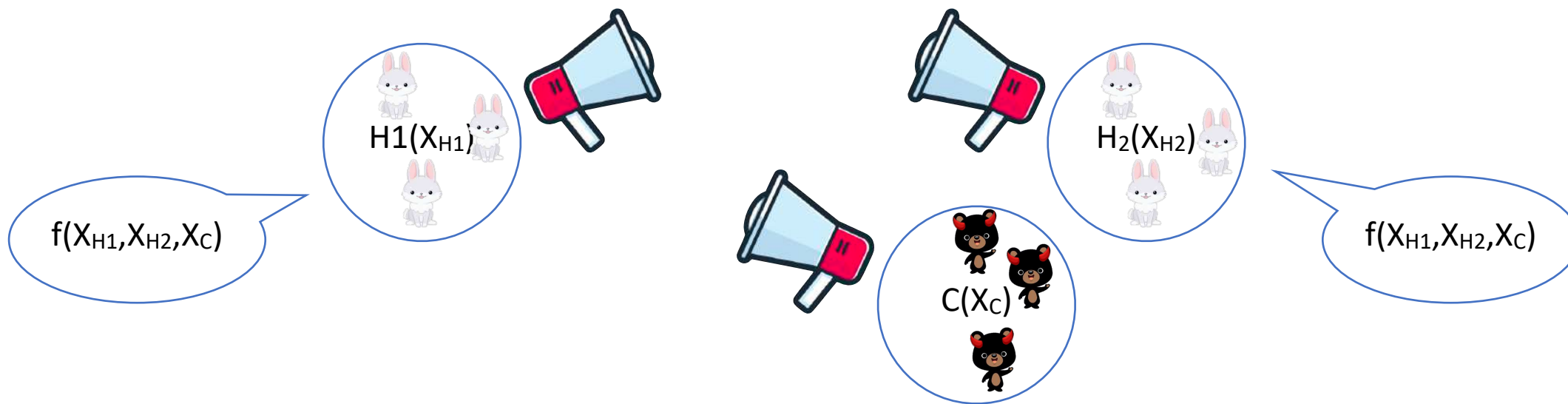
# $(t_d, t_m)$ -asyncP2P, BC

with private channels



round 1:  $(t_d, t_m)$ -asyncP2P

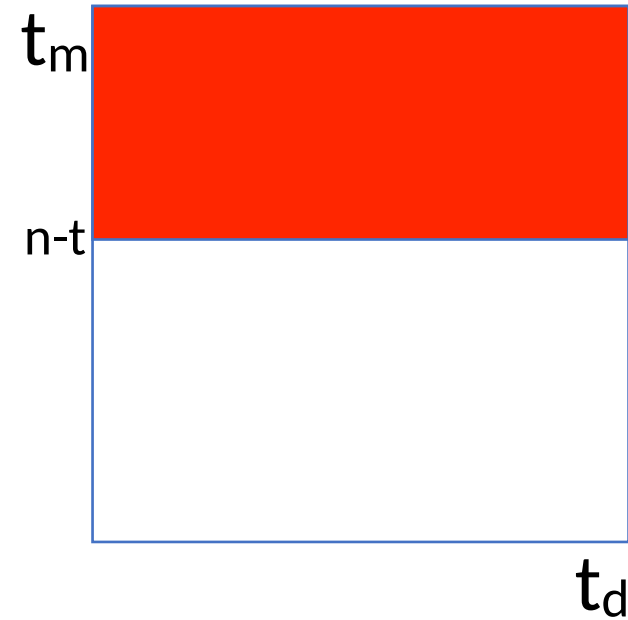
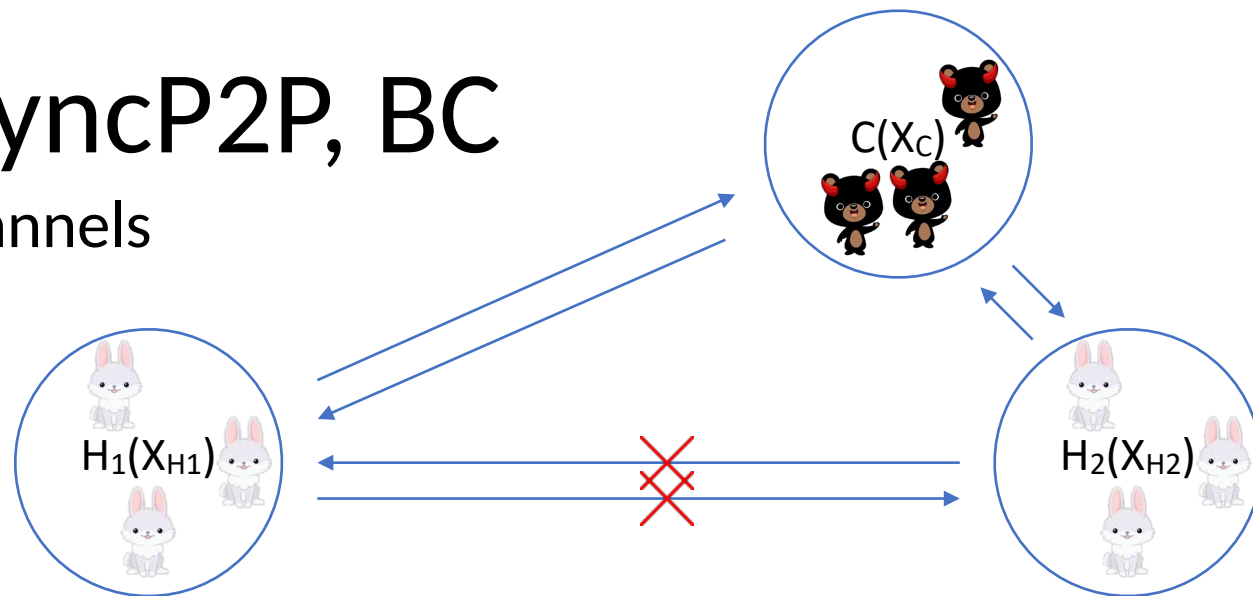
round 2: BC





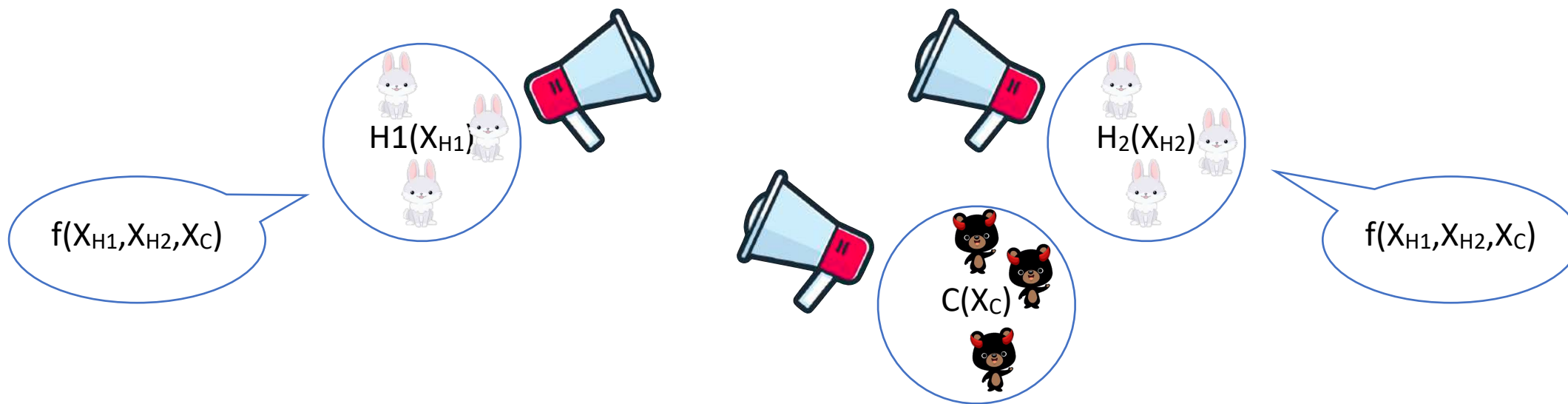
# $(t_d, t_m)$ -asyncP2P, BC

with private channels



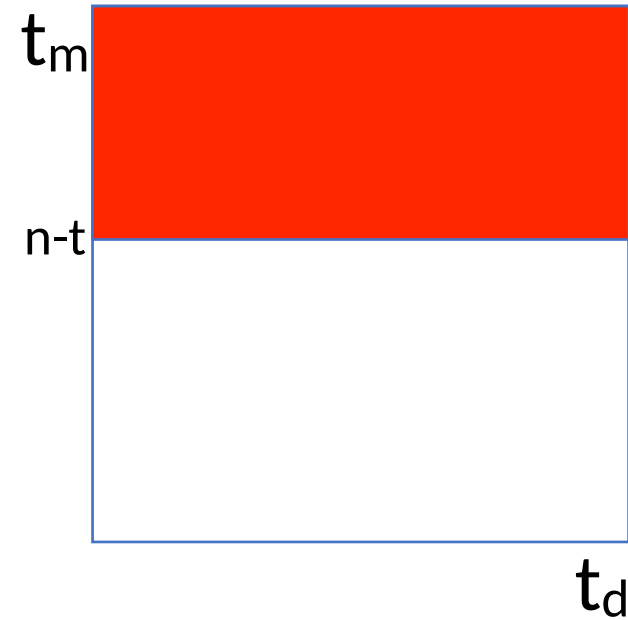
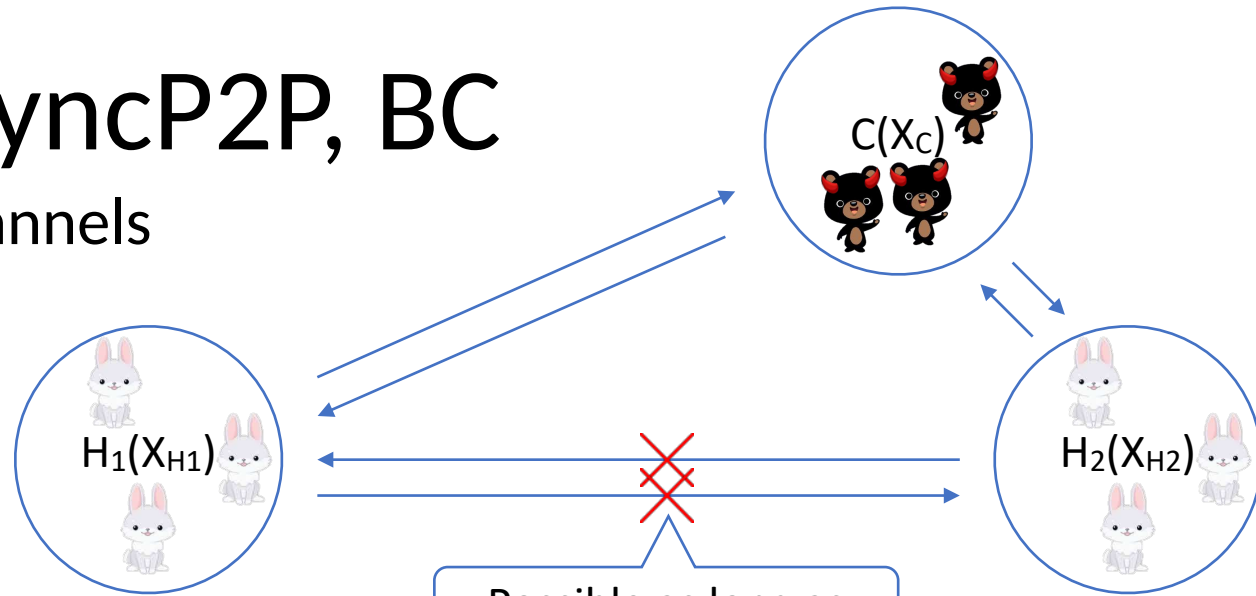
round 1:  $(t_d, t_m)$ -asyncP2P

round 2: BC



# $(t_d, t_m)$ -asyncP2P, BC

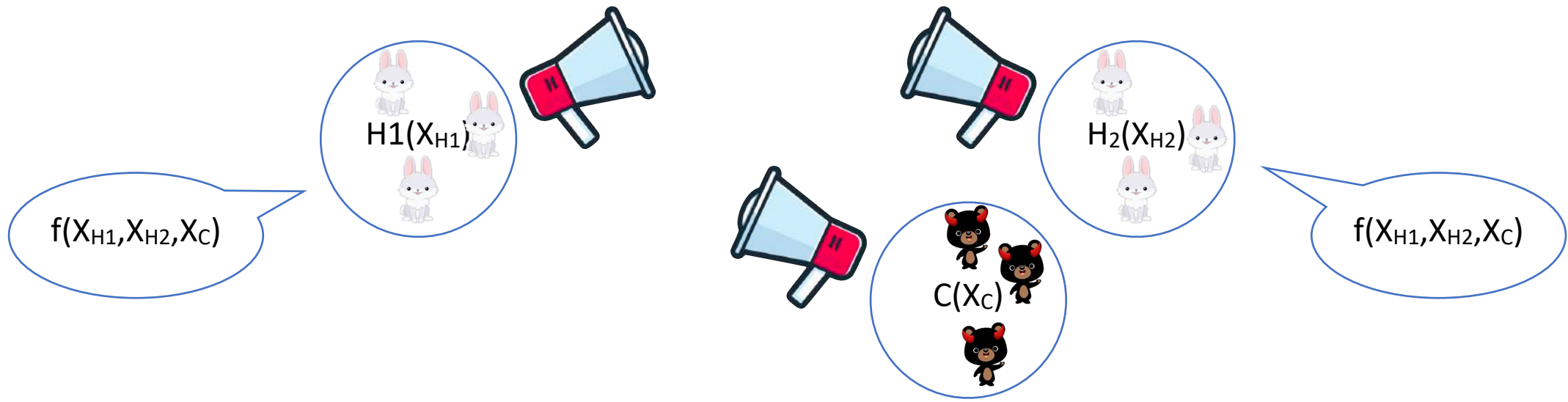
with private channels



Possible as long as  
 $t_m \geq (n-t)/2,$   
 $t_d \geq (n-t)/2$

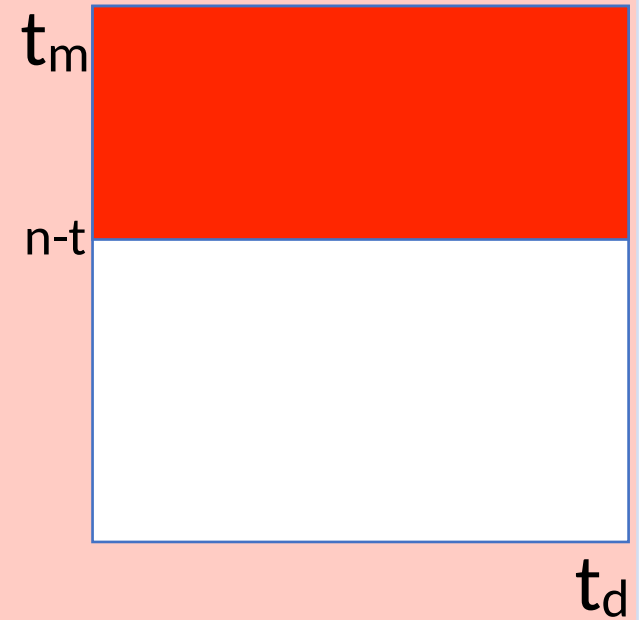
round 1:  $(t_d, t_m)$ -asyncP2P

round 2: BC



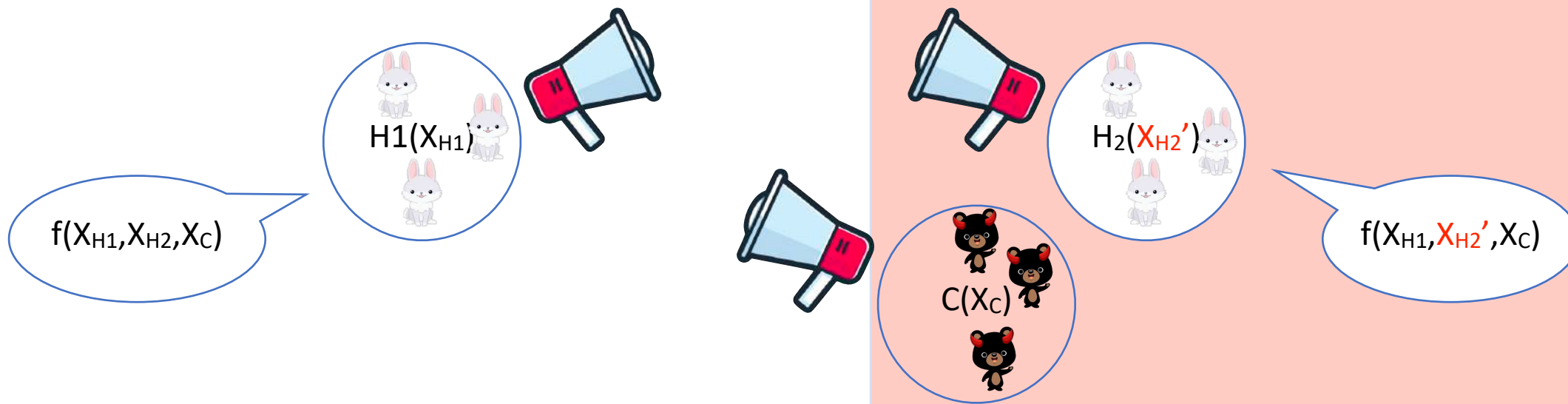
# $(t_d, t_m)$ -asyncP2P, BC

with private channels



round 1:  $(t_d, t_m)$ -asyncP2P

round 2: BC

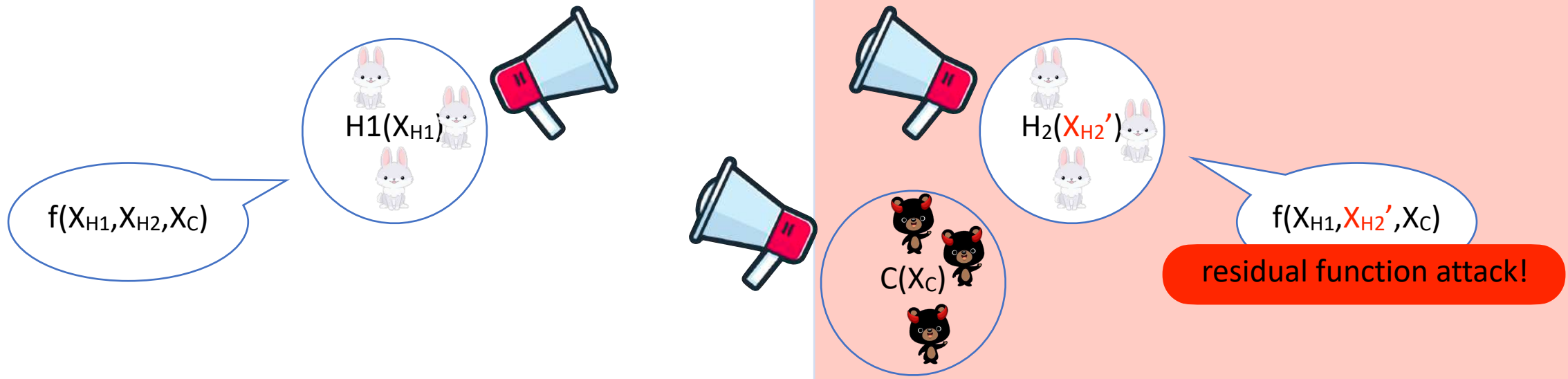
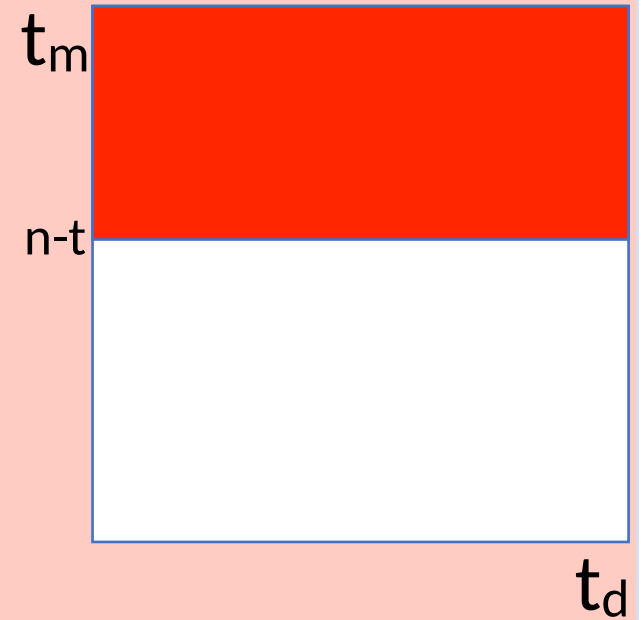


# $(t_d, t_m)$ -asyncP2P, BC

with private channels

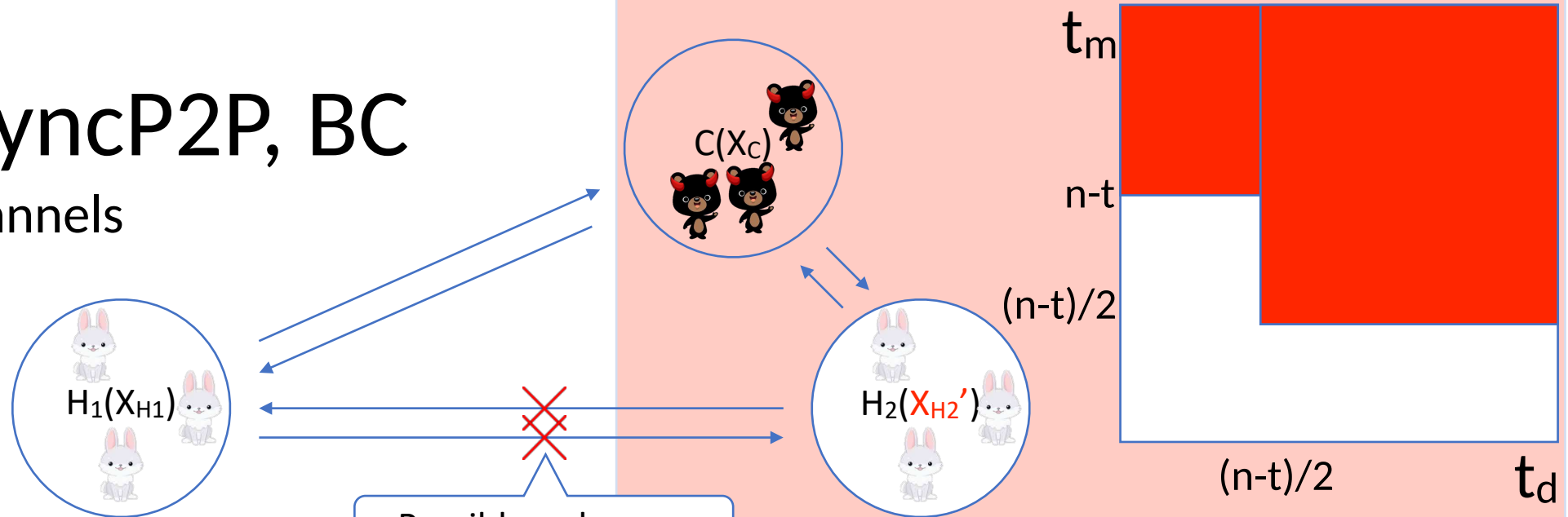


Possible as long as  
 $t_m \geq (n-t)/2$ ,  
 $t_d \geq (n-t)/2$



# $(t_d, t_m)$ -asyncP2P, BC

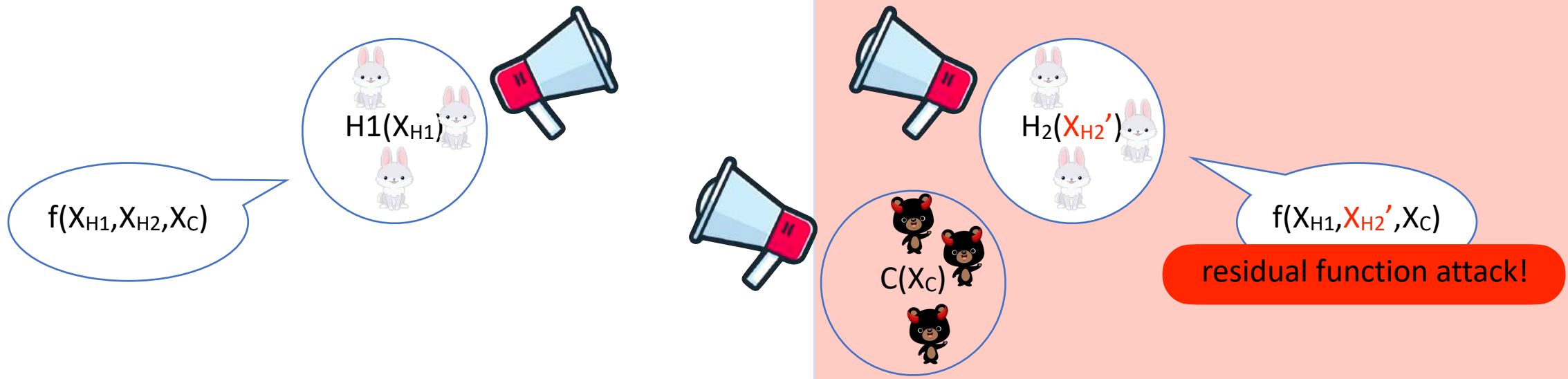
with private channels



Possible as long as  
 $t_m \geq (n-t)/2$ ,  
 $t_d \geq (n-t)/2$

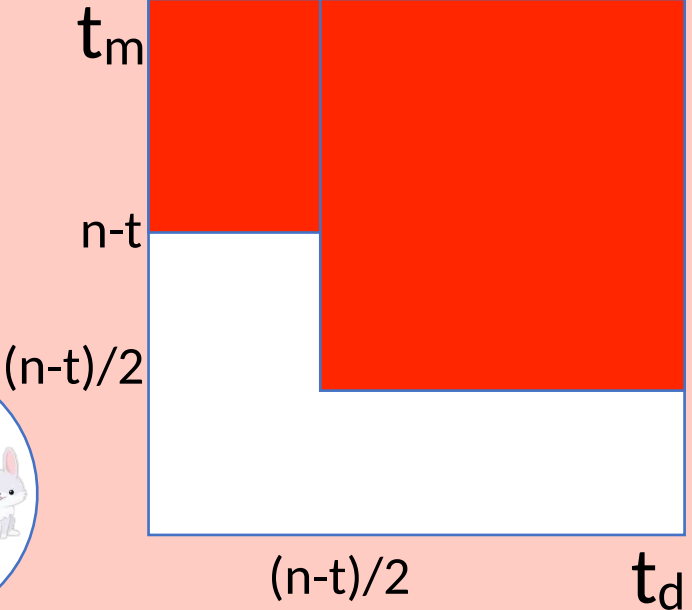
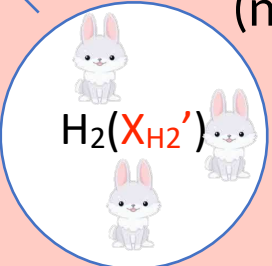
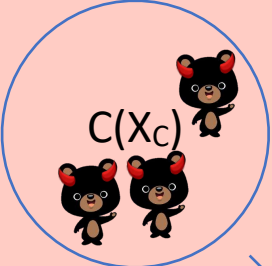
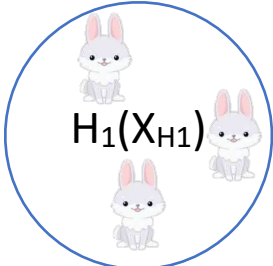
round 1:  $(t_d, t_m)$ -asyncP2P

round 2: BC



$f(X_{H1}, X_{H2}', X_c)$   
residual function attack!

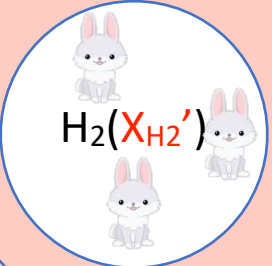
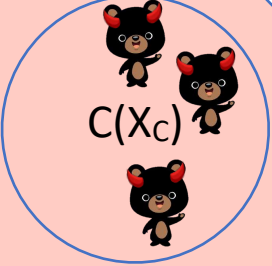
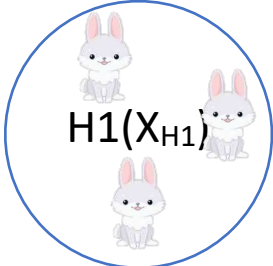
# $(t_d, t_m)$ -asyncP2P, BC



Possible as long as  
 $t_m \geq (n-t)/2$ ,  
 $t_d \geq (n-t)/2$

round 1:  $(t_d, t_m)$ -asyncP2P  
 round 2: BC

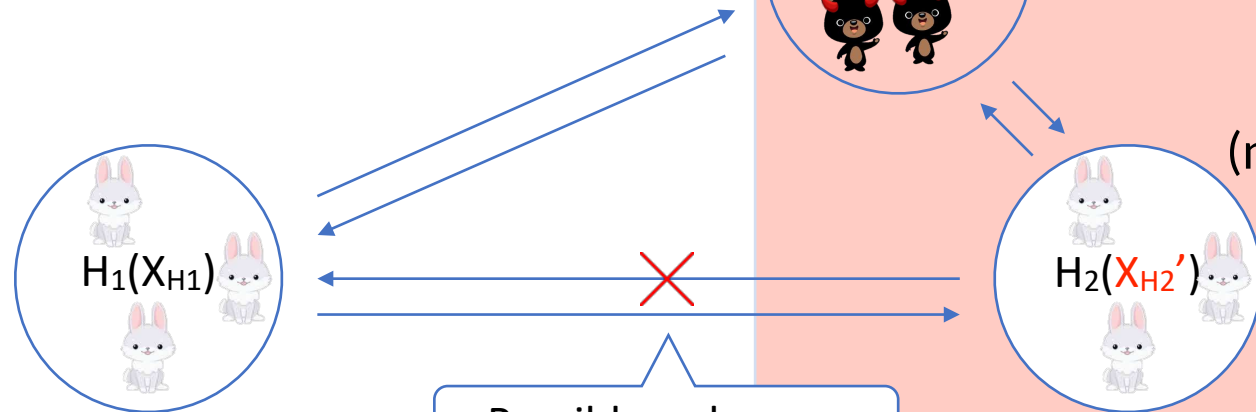
$f(X_{H1}, X_{H2}, X_C)$



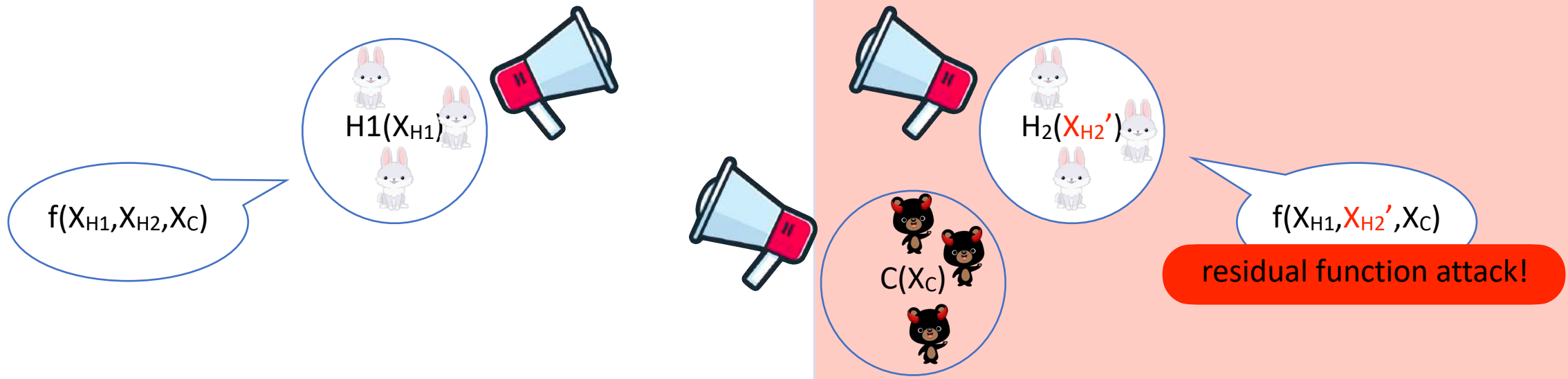
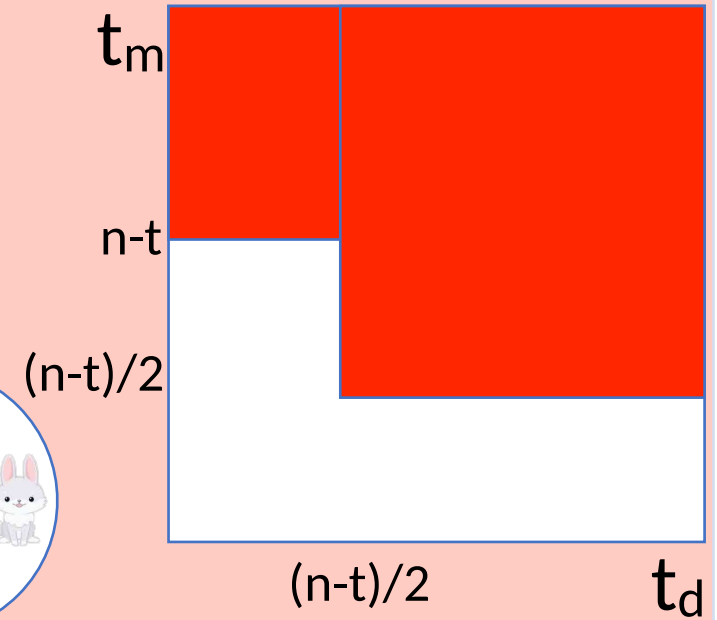
$f(X_{H1}, X_{H2}', X_C)$

residual function attack!

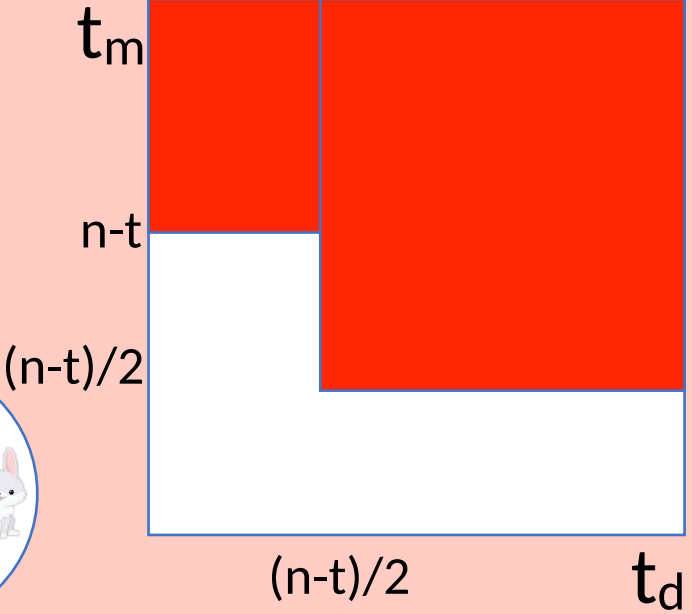
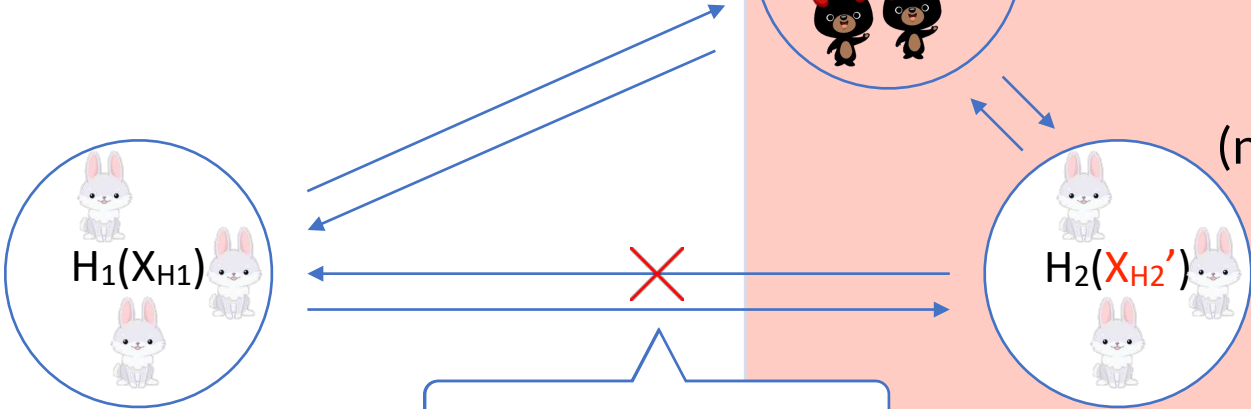
# $(t_d, t_m)$ -asyncP2P, BC



Possible as long as  
 $t_m \geq (n-t)/2$ ,  
 $t_d \geq (n-t)/2$

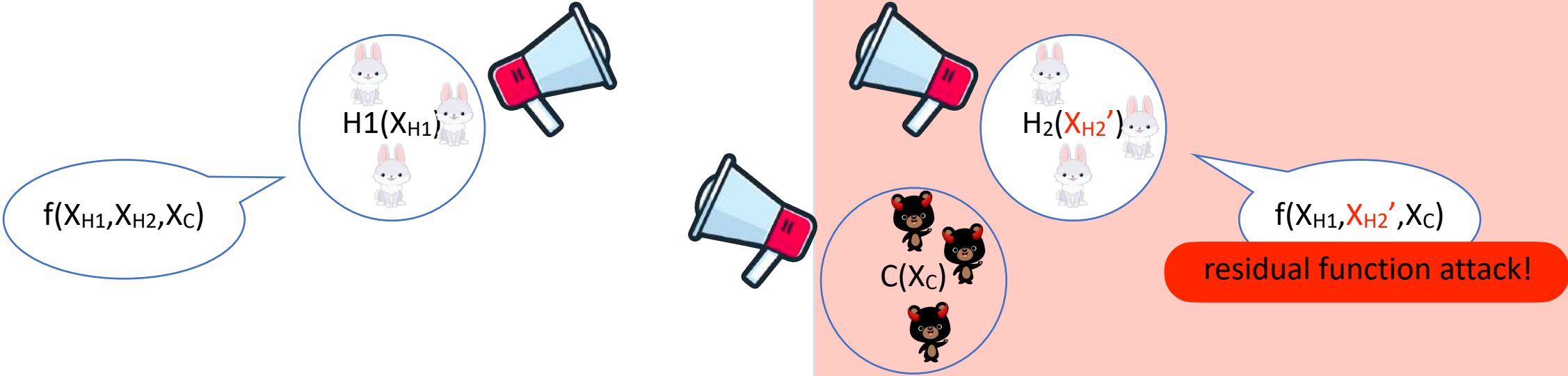


# $(t_d, t_m)$ -asyncP2P, BC



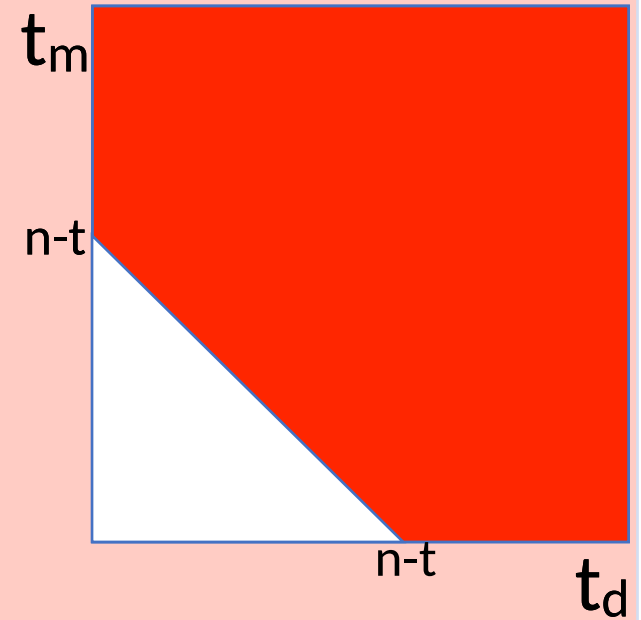
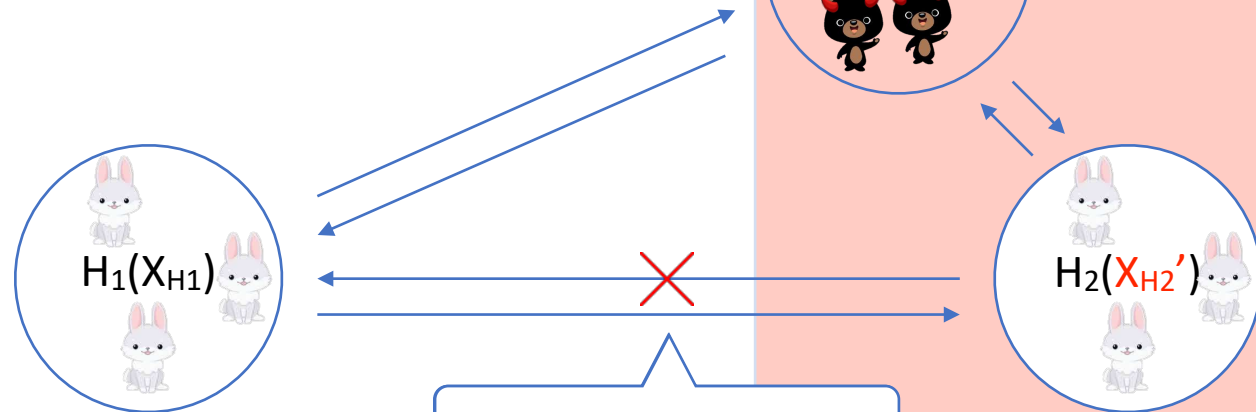
round 1:  $(t_d, t_m)$ -asyncP2P  
 -----  
 round 2: BC

Possible as long as  
 $t_m + t_d \geq n - t$



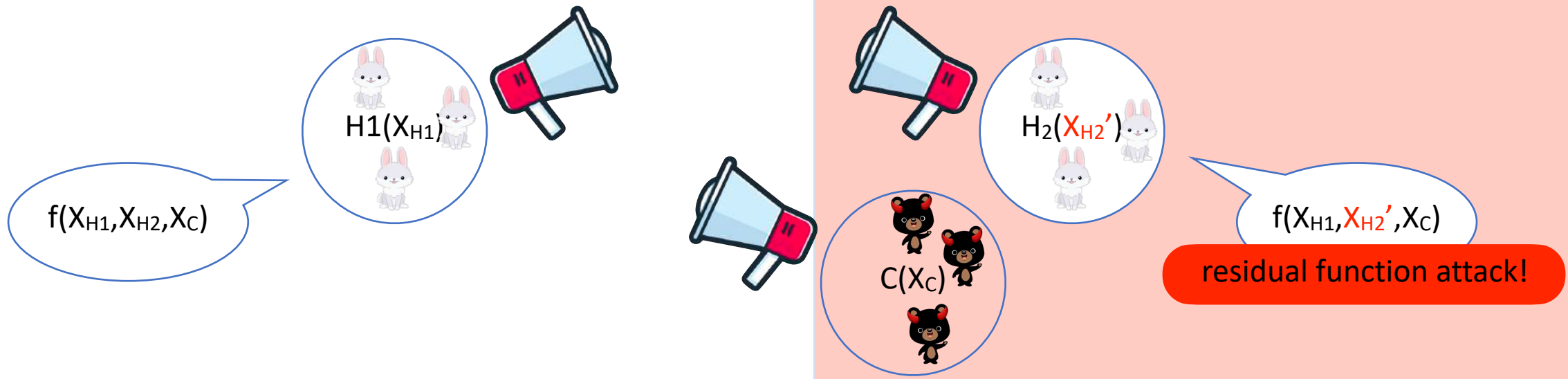


# $(t_d, t_m)$ -asyncP2P, BC

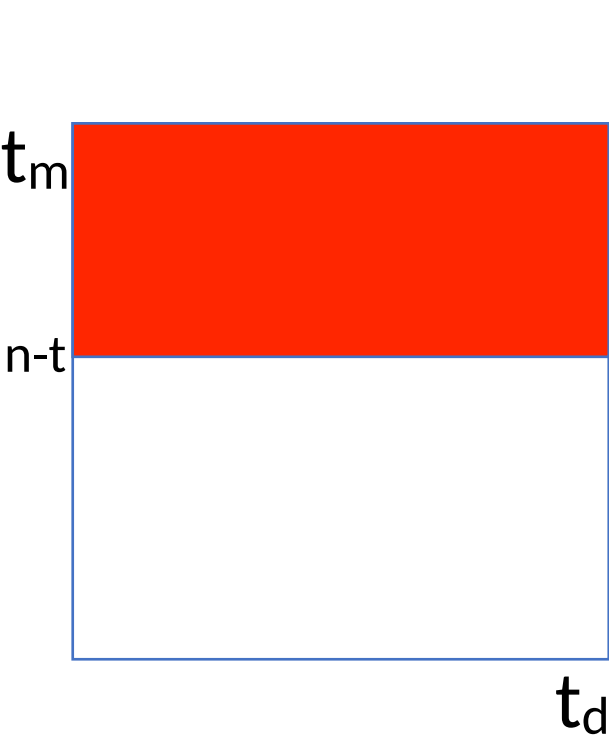


round 1:  $(t_d, t_m)$ -asyncP2P

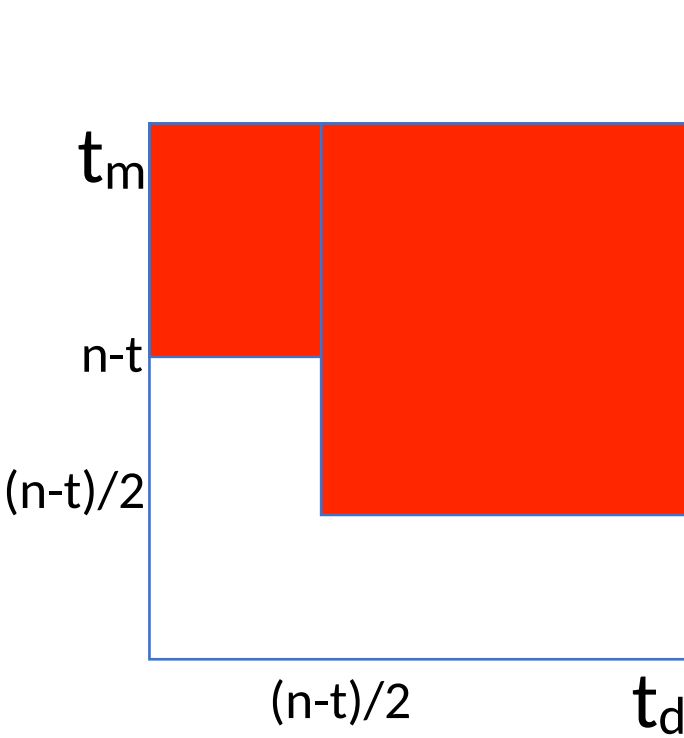
round 2: BC



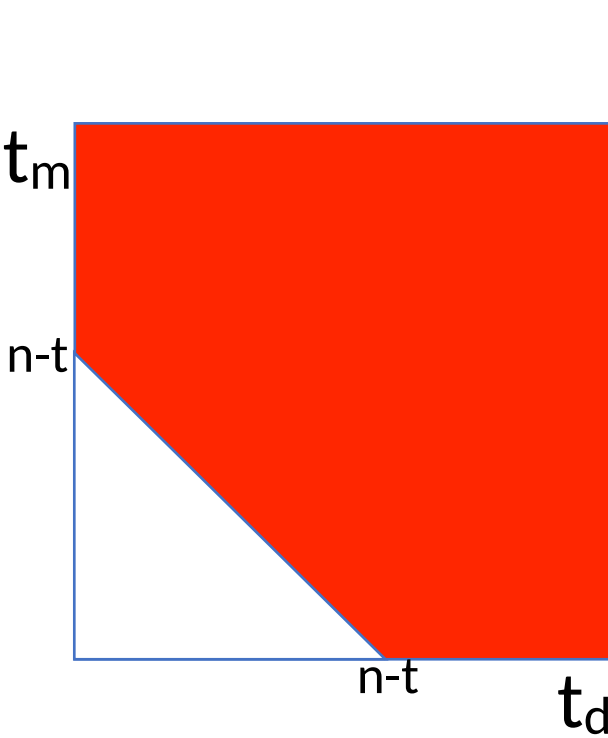
# $(t_d, t_m)$ -asyncP2P, BC



with PKI

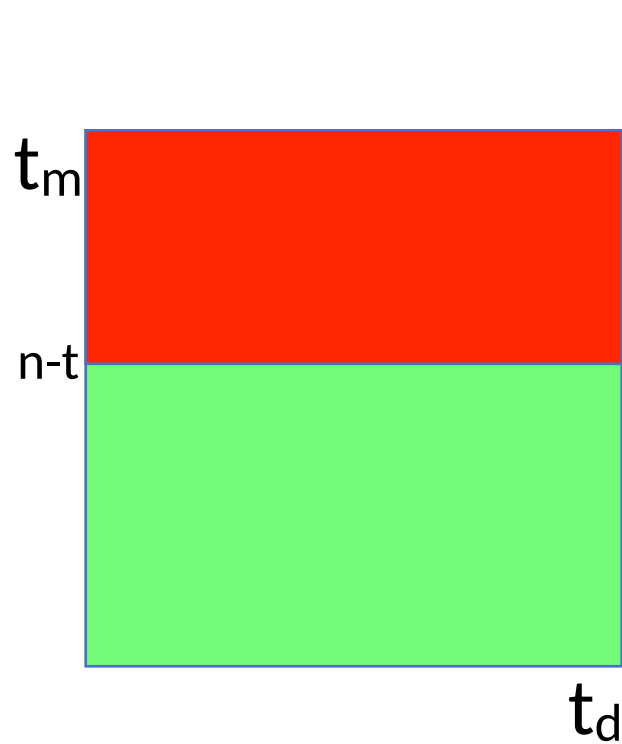


with private channels

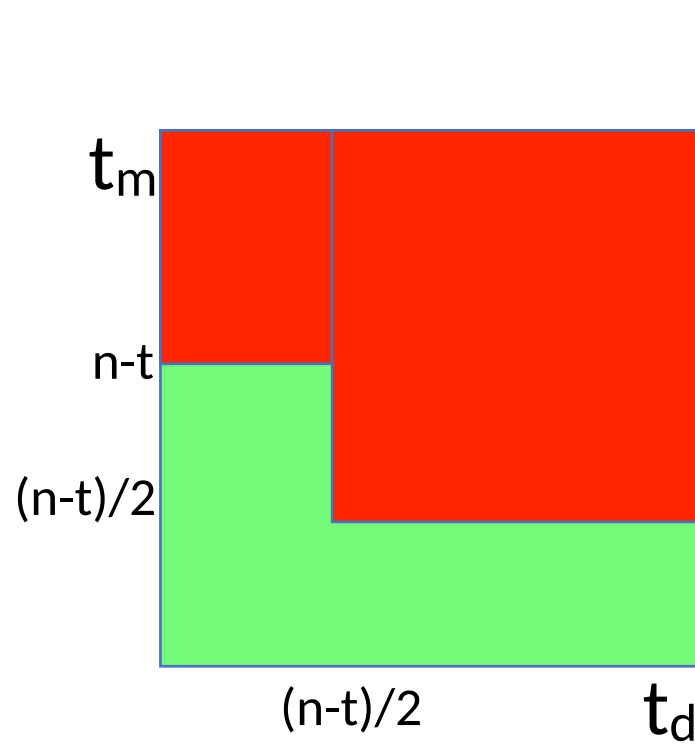


no PKI or private channels

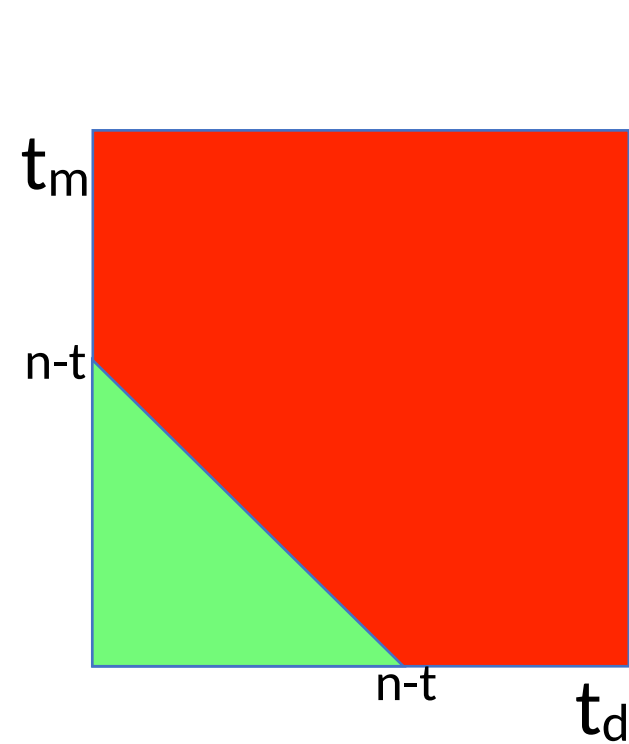
# $(t_d, t_m)$ -asyncP2P, BC



with PKI

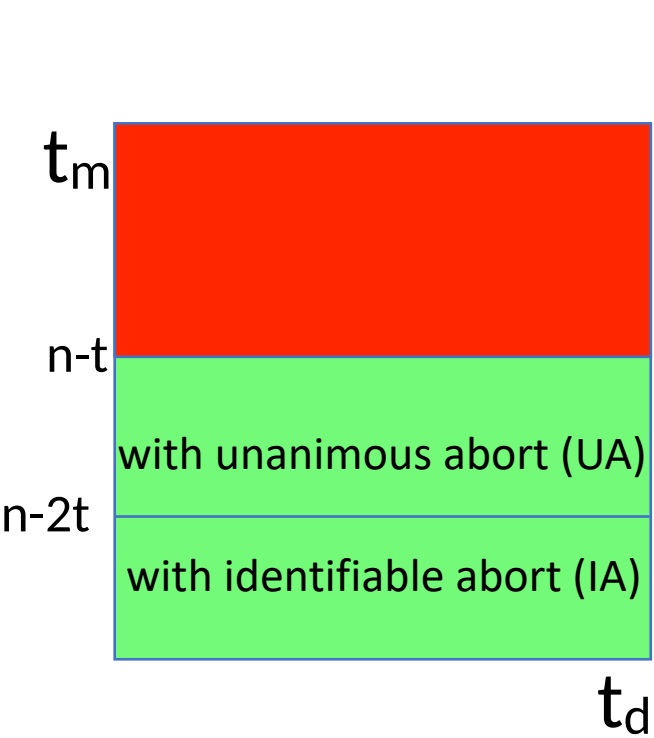


with private channels

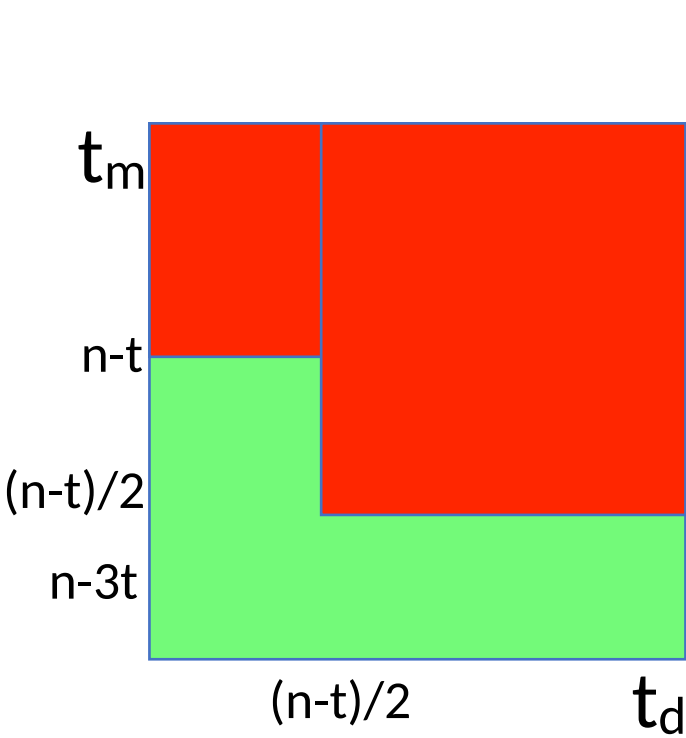


no PKI or private channels

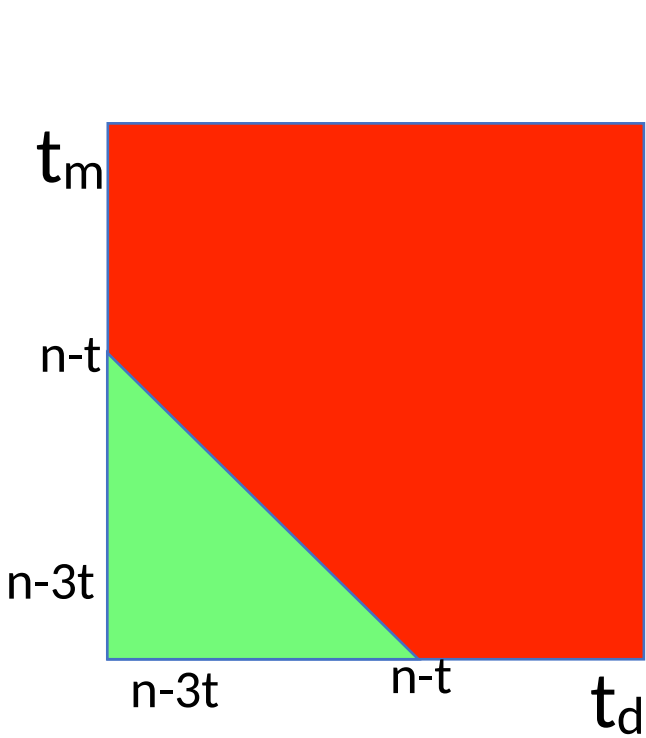
# $(t_d, t_m)$ -asyncP2P, BC



with PKI

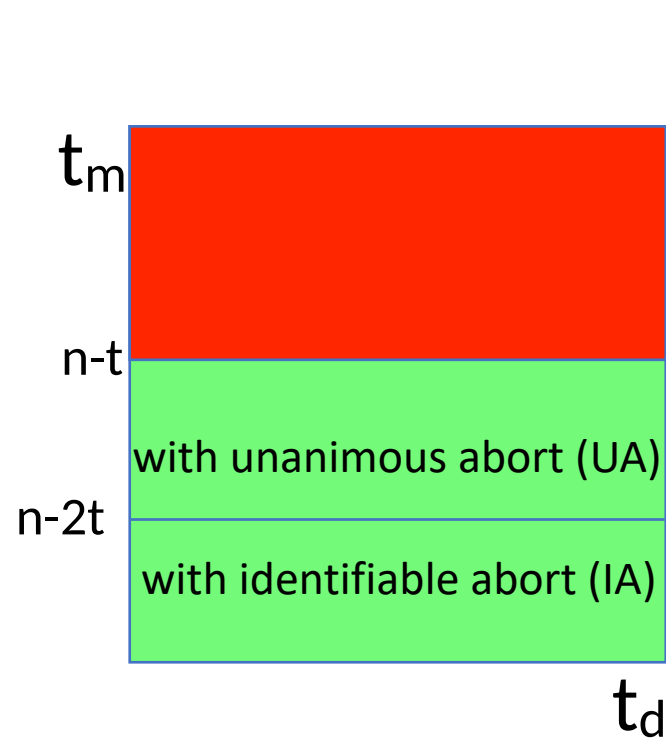


with private channels

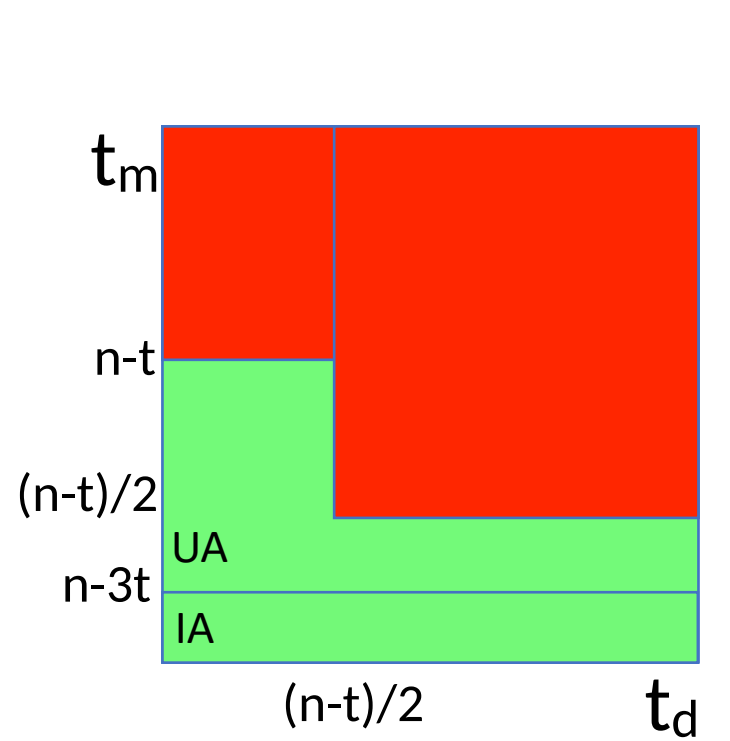


no PKI or private channels

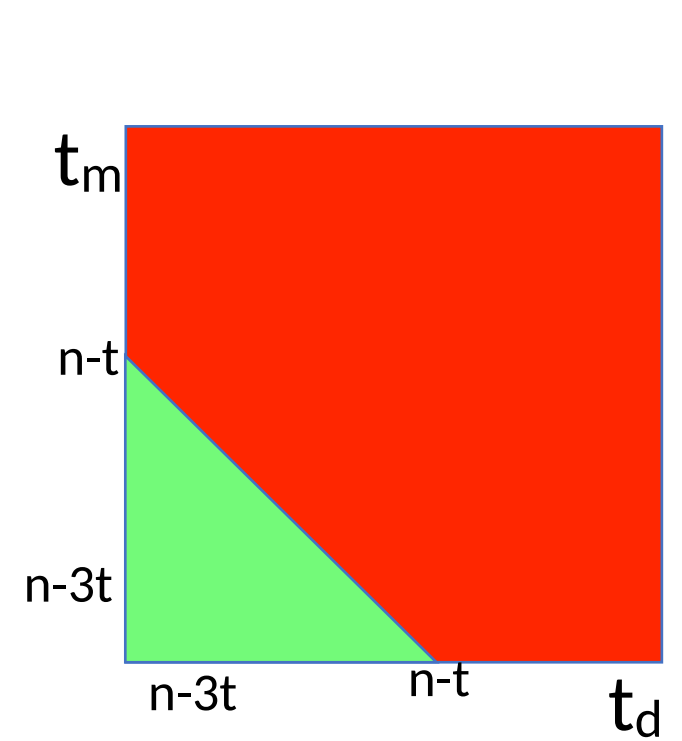
# $(t_d, t_m)$ -asyncP2P, BC



with PKI

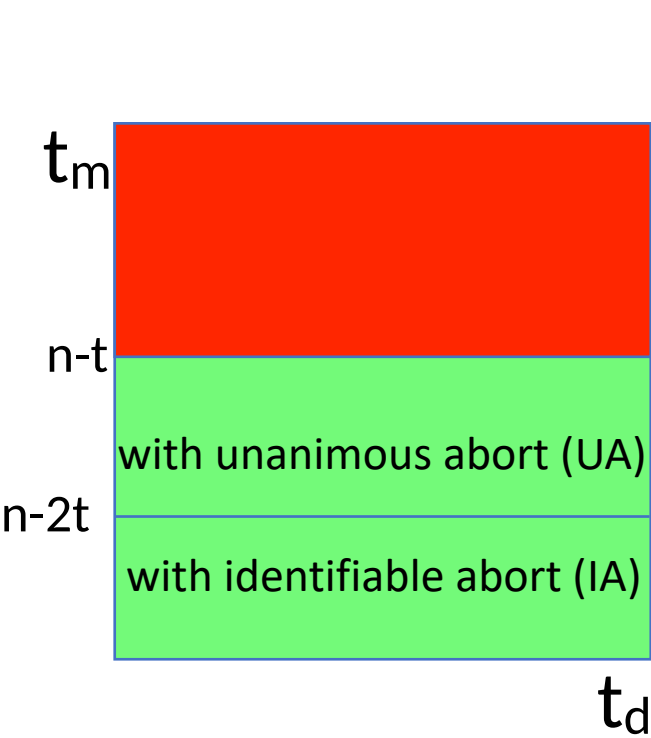


with private channels

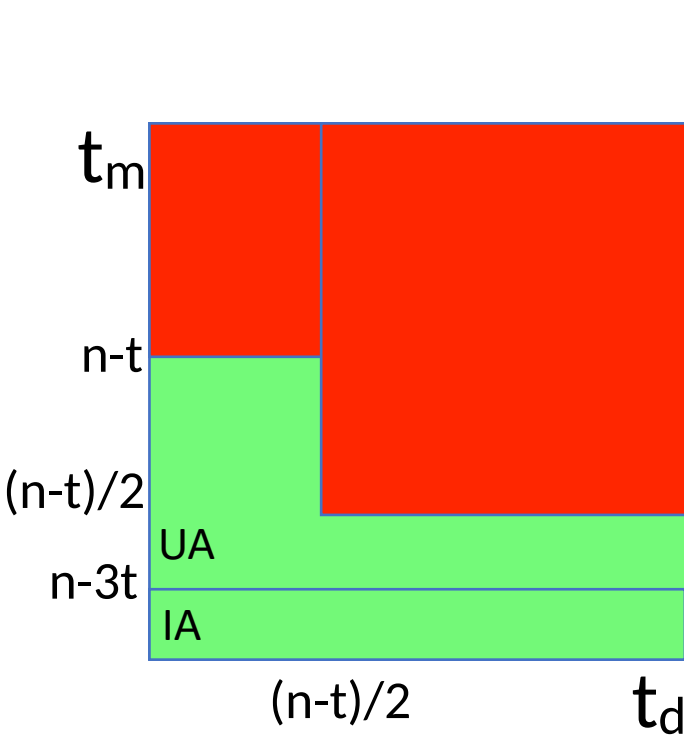


no PKI or private channels

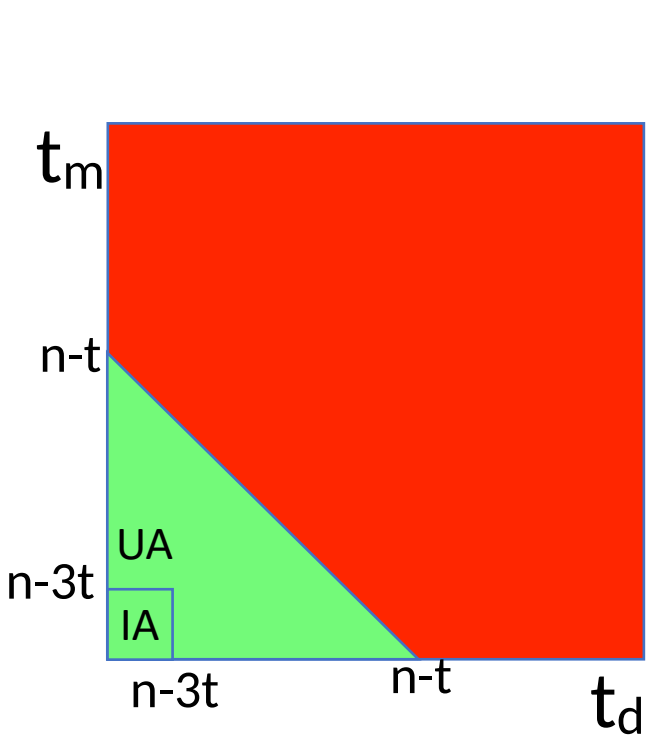
# $(t_d, t_m)$ -asyncP2P, BC



with PKI



with private channels



no PKI or private channels

# $(t_d, t_m)$ -asyncP2P, BC: Constructions

- Using tools from previous papers
  - Variants of one-or-nothing secret sharing
    - Do not support all values of  $t_d, t_m$
- New constructions from indistinguishability obfuscation
  - New primitive: puncturable sender-public key encryption
    - Inefficient / unrealistic building blocks

# Summary

- Our contributions:
  - New notion of  $(t_d, t_m)$ -asynchrony
  - Impossibility results
  - Constructions